

# Two-step MIR inequalities

Oktay Günlük

Math Sciences, IBM Research

joint work with Sanjeeb Dash and Marcos Goycoolea

- *Introduction*
- *Part 1:*
  - *Mixed-integer rounding / GMIC*
  - *Two-step MIR set*
  - *Two-step MIR procedure*
- *Part 2:*
  - *Gomory's Master Cyclic Group Polyhedron*
  - *Shooting experiment*
- *Part 3:*
  - *Separation*
  - *Computation*

# Cutting planes for Mixed-Integer Programs

*LP Relaxation:*

$$\min c^T x$$

$$Ax \geq b$$

*With Cuts:*

$$\min c^T x$$

$$Ax \geq b$$

$$\alpha_1^T x \geq d_1$$

*More Cuts:*

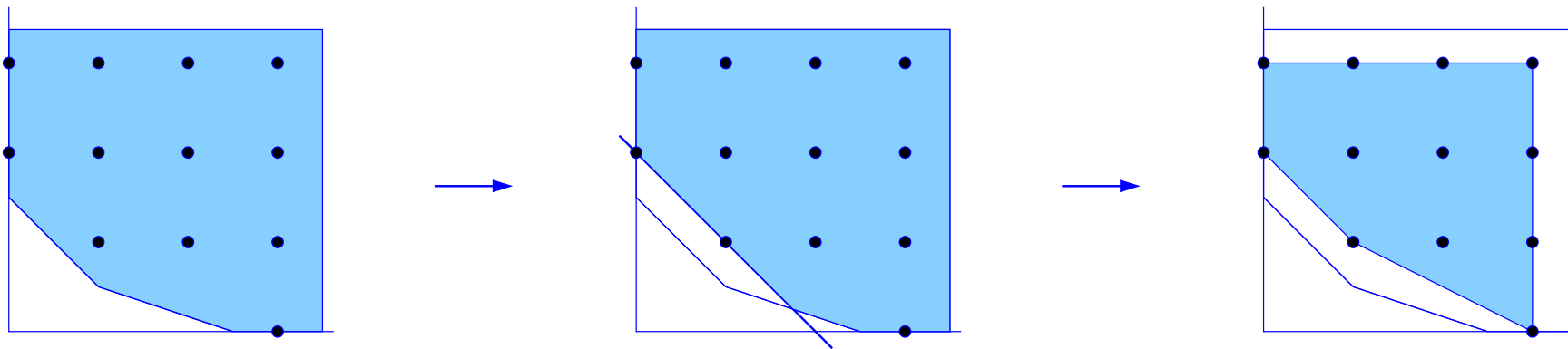
$$\min c^T x$$

$$Ax \geq b$$

$$\alpha_1^T x \geq d_1$$

$$\alpha_2^T x \geq d_2$$

:



# Single constraint relaxation

- Let

$$\sum_{j \in J} c_j v_j + \sum_{i \in I} a_i x_i \geq b$$

be an implied constraint of the MIP with  $v \in R_+^{|J|}$  and  $x \in Z_+^{|I|}$ .

- Consider the set

$$W' = \left\{ v \in R^2, x \in Z^{|I|} : v_1 - v_2 + \sum_{i \in I} a_i x_i \geq b, \quad x, v \geq 0 \right\}$$

Valid inequalities for  $W'$  are also valid for the mixed-integer program.

- Let  $\hat{b} = b - \lfloor b \rfloor$ , and  $\hat{a}_i = a_i - \lfloor a_i \rfloor$ . Define  $z = \sum_{j \in J} \lfloor a_j \rfloor x_j - \lfloor b \rfloor$ .  
Define the relaxed set:

$$W = \left\{ v \in R^2, x \in Z^{|I|}, z \in Z : v_1 - v_2 + \sum_{i \in I} \hat{a}_i x_i + z \geq \hat{b}, \quad x, v \geq 0 \right\}$$

- Goal: derive cuts for MIP using valid inequalities for  $W$

# Basic mixed-integer rounding set (Wolsey '98)

Let

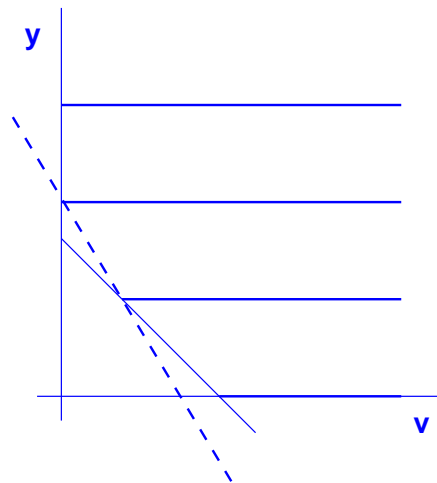
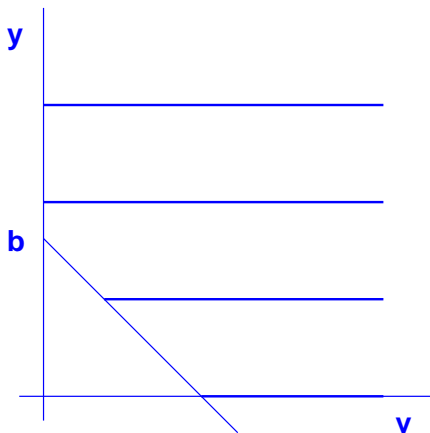
$$Q = \{v \in \mathbb{R}, y \in \mathbb{Z} : v + y \geq b, v \geq 0\}$$

then, MIR Inequality:

$$v \geq \hat{b}([\bar{b}] - y)$$

where  $\hat{b} = b - \lfloor b \rfloor$ , is valid for  $Q$  and

$$\text{conv}(Q) = \{v, y \in \mathbb{R} : v + y \geq b, v + \hat{b}y \geq \hat{b} \lfloor b \rfloor, v \geq 0\}.$$



# MIR inequalities for $W$

Base System:

$$v_1 - v_2 + \sum_{\hat{a}_i < \hat{b}} \hat{a}_i x_i + \sum_{\hat{a}_i \geq \hat{b}} \hat{a}_i x_i + z \geq \hat{b}.$$

Relax:

$$v_1 + \sum_{\hat{a}_i < \hat{b}} \hat{a}_i x_i + \sum_{\hat{a}_i \geq \hat{b}} x_i + z \geq \hat{b}.$$

Apply MIR:

$$v_1 + \sum_{\hat{a}_i < \hat{b}} \hat{a}_i x_i + \hat{b} \sum_{\hat{a}_i \geq \hat{b}} x_i + z \geq \hat{b}.$$

- If the base inequality comes from a row of the simplex tableau:

*MIR inequality = Gomory mixed-integer cut.*

*(Marchand and Wolsey '98)*

## Another simple set

- Let

$$Q^2 = \left\{ v \in R, y, z \in Z : v + \alpha y + z \geq \beta, v, y, z \geq 0 \right\},$$

where  $0 < \alpha < \beta < 1$ .

- The complete description of  $Q^2$  is:

$$\text{conv}(Q^2) = \left\{ v, y, z \in R : \begin{array}{l} v + \alpha y + \beta z \geq \beta \\ \frac{v}{\beta - \alpha} + y + \lceil \beta/\alpha \rceil z \geq \lceil \beta/\alpha \rceil \\ v, y, z \geq 0 \end{array} \right\}$$

where the second inequality is called the two-step MIR inequality (DG'03)

## Two-step MIR

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*Original:*

$$(v + \alpha y) + z \geq \beta.$$

*MIR with  $z$  integral:*

$$(v + \alpha y) + \beta z \geq \beta.$$

*Divide by  $\alpha$ :*

$$\frac{v}{\alpha} + y + \frac{\beta}{\alpha} z \geq \frac{\beta}{\alpha}.$$

*Relax coefficient of  $z$ :*

$$\frac{v}{\alpha} + (y + \lceil \frac{\beta}{\alpha} \rceil z) \geq \frac{\beta}{\alpha}.$$

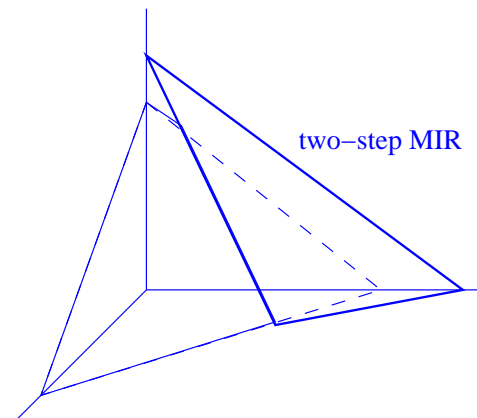
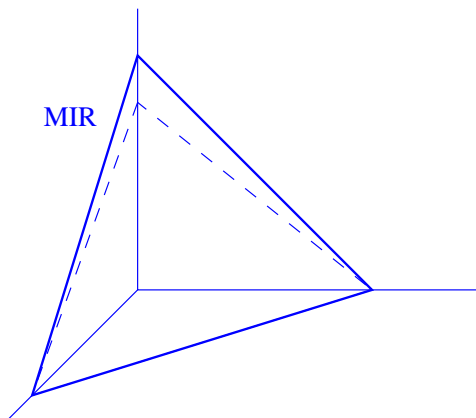
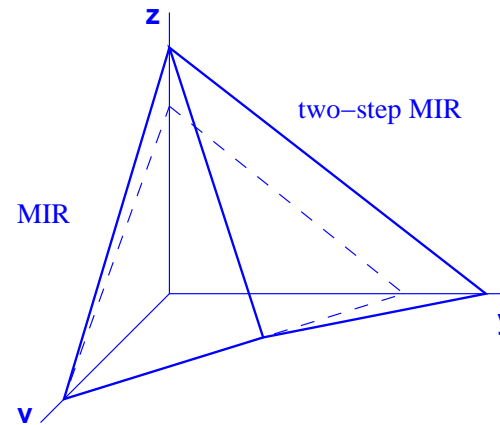
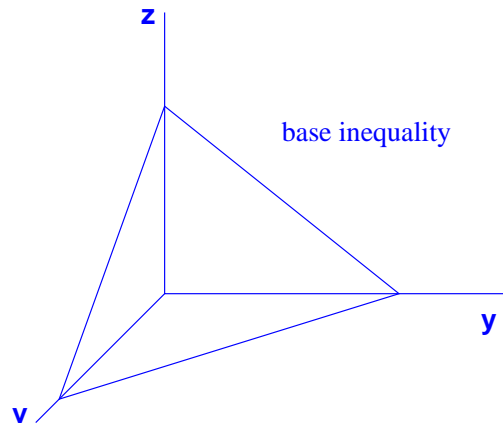
*MIR with  $v/\alpha$  continuous:*

$$\frac{v}{\beta - \alpha \lceil \beta/\alpha \rceil} + y + \lceil \beta/\alpha \rceil z \geq \lceil \beta/\alpha \rceil$$



# Facets of $Q^2$

$$Q^2 = \{v \in R, y, z \in Z : v + \alpha y + z \geq \beta, v, y, z \geq 0\}$$



## A slightly more general set

What if  $z \geq 0$  does not hold?

$$Q^{2R} = \left\{ v \in R, y, z \in Z : v + \alpha y + z \geq \beta, v, y \geq 0 \right\}.$$

- This set is more useful for generating cuts.

**Theorem** (DG'03) If  $1/\alpha \geq \lceil \beta/\alpha \rceil$ , then the two-step MIR inequality

$$\frac{v}{\beta - \alpha \lfloor \beta/\alpha \rfloor} + y + \lceil \beta/\alpha \rceil z \geq \lceil \beta/\alpha \rceil$$

is valid for  $Q^{2R}$ . Furthermore, it is facet defining provided that  $\beta/\alpha \notin Z$ .

- Not a complete description of  $Q^{2R}$ .
- Agra and Constantino (2003) and Atamtürk and Rajan (2003) provide a polynomial algorithm to enumerate all facets of  $Q^{2R}$ .

## Derivation of two-step MIR for $Q^{2R}$

Original:

$$(v + \alpha y) + z \geq \beta.$$

Applying MIR with  $z$  integral:

$$(v + \alpha y) + \beta z \geq \beta.$$

Therefore

$$\frac{v}{\alpha} + y + \frac{\beta}{\alpha}z \geq \frac{\beta}{\alpha} \quad \text{and} \quad \frac{v}{\alpha} + y + \frac{1}{\alpha}z \geq \frac{\beta}{\alpha}$$

are valid for  $Q^{2R}$ .

When  $1/\alpha \geq \lceil \beta/\alpha \rceil$ , convex combination gives:

$$\frac{v}{\alpha} + (y + \lceil \frac{\beta}{\alpha} \rceil z) \geq \frac{\beta}{\alpha}$$

Applying MIR with  $(y + \lceil \beta/\alpha \rceil z)$  integral gives the two-step MIR inequality.

# Example

Base:

$$v + 0.4y + z \geq 1.7$$

MIR with  $z$  integral:

$$(v + 0.4y) + 0.7z \geq 1.4$$

divide by 0.4:

$$2.5v + y + 1.75z \geq 3.5$$

Relaxation

Positive combination with base

$$2.5v + (y + 2z) \geq 3.5$$

$$2.5v + y + 1.75(z - 1) \geq 1.75$$

$$2.5v + y + 2(z - 1) \geq 1.75$$

$$2.5v + (y + 2z) \geq 3.75$$

MIR with  $(y + 2z)$  integral:

$$2.5v + 0.5(y + 2z) \geq 2$$

$$2.5v + 0.75(y + 2z) \geq 3$$

## Applying two-step MIR's to $W$

We know that if  $v + \alpha y + z \geq \beta$  is valid, and  $v, y \geq 0$ , then

$$\frac{v}{\beta - \alpha \lfloor \beta/\alpha \rfloor} + y + \lceil \beta/\alpha \rceil z \geq \lceil \beta/\alpha \rceil$$

is also valid if  $1/\alpha \geq \lceil \beta/\alpha \rceil$ .

To apply to  $v_1 - v_2 + \sum_{i \in I} \hat{a}_i x_i + z \geq \hat{b}$ , where  $v, x \geq 0$  and  $x, z$  integer,

- Chose an  $\alpha$  such that  $1/\alpha \geq \lceil \hat{b}/\alpha \rceil > \hat{b}/\alpha$
- Relax the inequality to obtain an inequality of the form

$$\underbrace{(v_1 + \dots)}_{\geq 0} + \alpha \underbrace{(\dots)}_{\in \mathbb{Z}_+} + \underbrace{(\dots + z)}_{\in \mathbb{Z}} \geq \hat{b}$$

- Apply two-step MIR inequality to the relaxed inequality

## Applying two-step MIR's – cont'd

- Depending on  $\hat{a}_i$ , we can rewrite/relax it as

- $\hat{a}_i \rightarrow 1$ , or,

- $\hat{a}_i \rightarrow \alpha(\lceil \hat{a}_i/\alpha \rceil)$ , or,

- $\hat{a}_i = \alpha(\lfloor \hat{a}_i/\alpha \rfloor) + (\hat{a}_i - \alpha \lfloor \hat{a}_i/\alpha \rfloor)$

- Partitioning  $I$  into  $I_1, I_2$  and  $I_3$  in a certain way gives the “best” relaxed inequality

$$v_1 + \sum_{j \in I_1} (\hat{a}_j - \alpha \lfloor \hat{a}_j/\alpha \rfloor) x_j + \alpha \sum_{j \in I_1} \lfloor \hat{a}_j/\alpha \rfloor x_j + \sum_{j \in I_2} \lceil \hat{a}_j/\alpha \rceil x_j + \sum_{j \in I_3} x_j + z \geq \hat{b}.$$

- Applied to  $v_1 - v_2 + \sum_{i \in I} a_i x_i \geq b$  this leads to the two-step MIR inequality

$$v_1 + \sum_{j \in I} \rho \tau \lfloor a_j \rfloor + \min\{\rho \tau, k_j \rho + \hat{a}_j - k_j \alpha, l_j \rho\} x_j \geq \rho \tau \lceil b \rceil,$$

where  $\tau = \lceil \hat{b}/\alpha \rceil$ ,  $\rho = \hat{b} - \alpha \lfloor \hat{b}/\alpha \rfloor$ ,  $k_i = \lfloor \hat{a}_i/\alpha \rfloor$  and  $l_i = \lceil \hat{a}_i/\alpha \rceil$ .

# Gomory's Master Cyclic Group Polyhedron

Gomory's Master Cyclic Group Polyhedron for  $n, r \in \mathbb{Z}$  and  $n > r > 0$ :

$$P(n, r) = \text{conv}\{w \in \mathbb{Z}_+^{n-1} : \sum_{i=1}^{n-1} i w_i = r \pmod{n}\}$$

or, equivalently,

$$P(n, r) = \text{conv}\left\{w \in \mathbb{Z}_+^{n-1} : \exists z \in \mathbb{Z} \text{ s.t. } \sum_{i=1}^{n-1} \binom{i}{n} w_i - z = \frac{r}{n}\right\}$$

Introduced by Gomory (1969). Also studied by GJ '72, GJE '03, AGJE '03.

For rational data,  $W$ , without the continuous variables, looks like a restriction of  $P(n, r)$

$$Y = \left\{x \in \mathbb{Z}_+^{|J|}, z \in \mathbb{Z} : \sum_{j \in J} \hat{a}_j x_j + z = \hat{b}\right\}$$

For some large  $\bar{n}$ ,  $Y$  is a face of  $P(\bar{n}, \bar{n} \cdot \hat{b})$ .

## Sub-additive functions on $(0, 1)$

- For all facets of  $P(n, r)$ , there is a sub-additive function that generates it.
- Ex: Gomory mixed-integer cut on  $(\sum_{j \in J} \hat{a}_j x_j + z = \hat{b})$ :

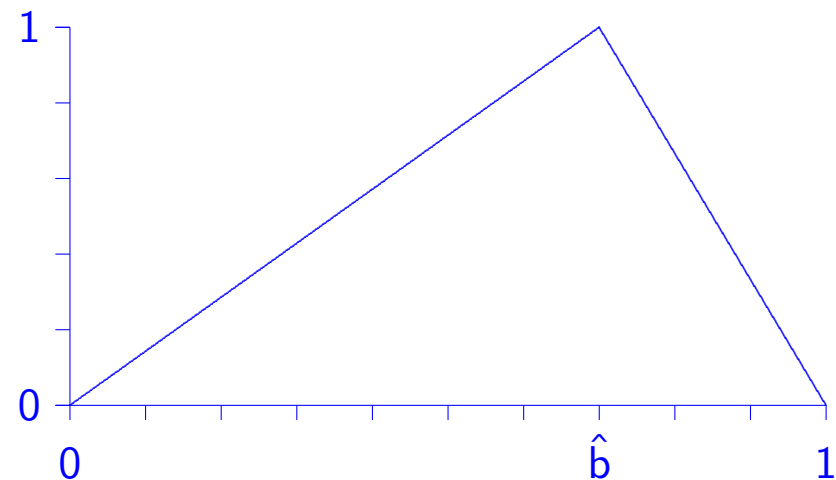
$$\sum_{\hat{a}_j < \hat{b}} \frac{\hat{a}_j}{\hat{b}} x_j + \sum_{\hat{a}_j \geq \hat{b}} \frac{1 - \hat{a}_j}{1 - \hat{b}} x_j \geq 1.$$

can also be written as

$$\sum_{j \in J} f^{\hat{b}}(\hat{a}_j) x_j \geq 1$$

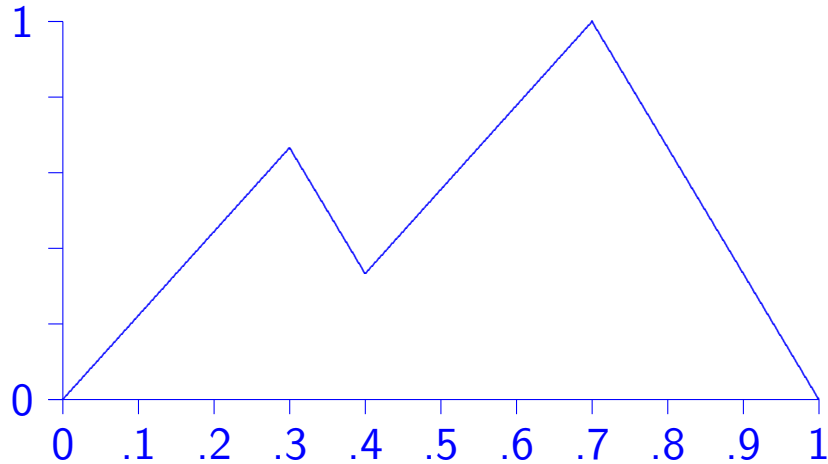
where

$$f^{\hat{b}}(\hat{v}) = \begin{cases} \hat{v}/\hat{b} & \text{if } \hat{v} < \hat{b}, \\ (1 - \hat{v})/(1 - \hat{b}) & \text{if } \hat{v} \geq \hat{b}, \end{cases}$$

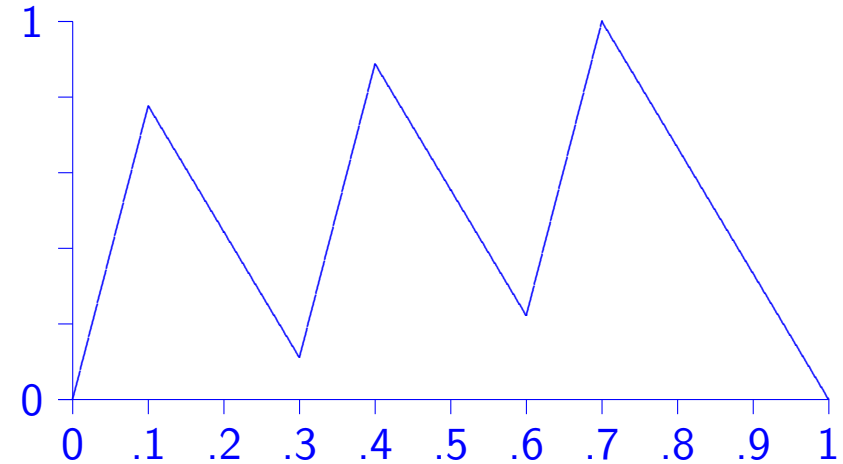




## Two-step MIR functions



$$\hat{b} = .7, \alpha = .4, \lceil \hat{b}/\alpha \rceil = 2$$



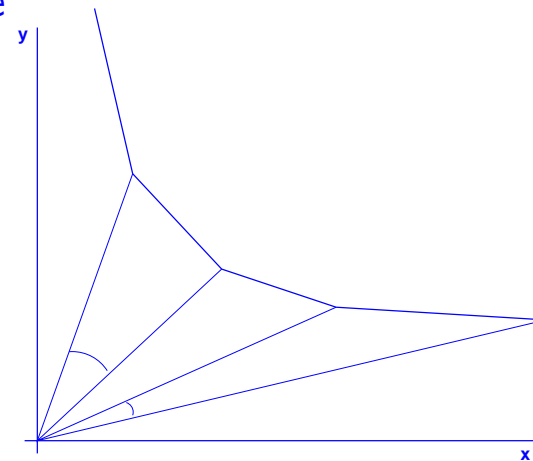
$$\hat{b} = .7, \alpha = .3, \lceil \hat{b}/\alpha \rceil = 3$$

- The two-step MIR function defines a facet of  $P(n, r)$  provided that  $\alpha$  is a multiple of  $1/n$  and  $1/\alpha > \lceil \hat{b}/\alpha \rceil$ . (DG'03)
- Scaled two-step MIR functions also define a facet of  $P(n, r)$ . (DG'03)
- These inequalities generalize the 2slope inequalities of Aroaz, Evans, Gomory, and Johnson (2003)
- Gomory and Johnson (1972) call these functions 2-slope functions.

# Gomory's Shooting Experiment

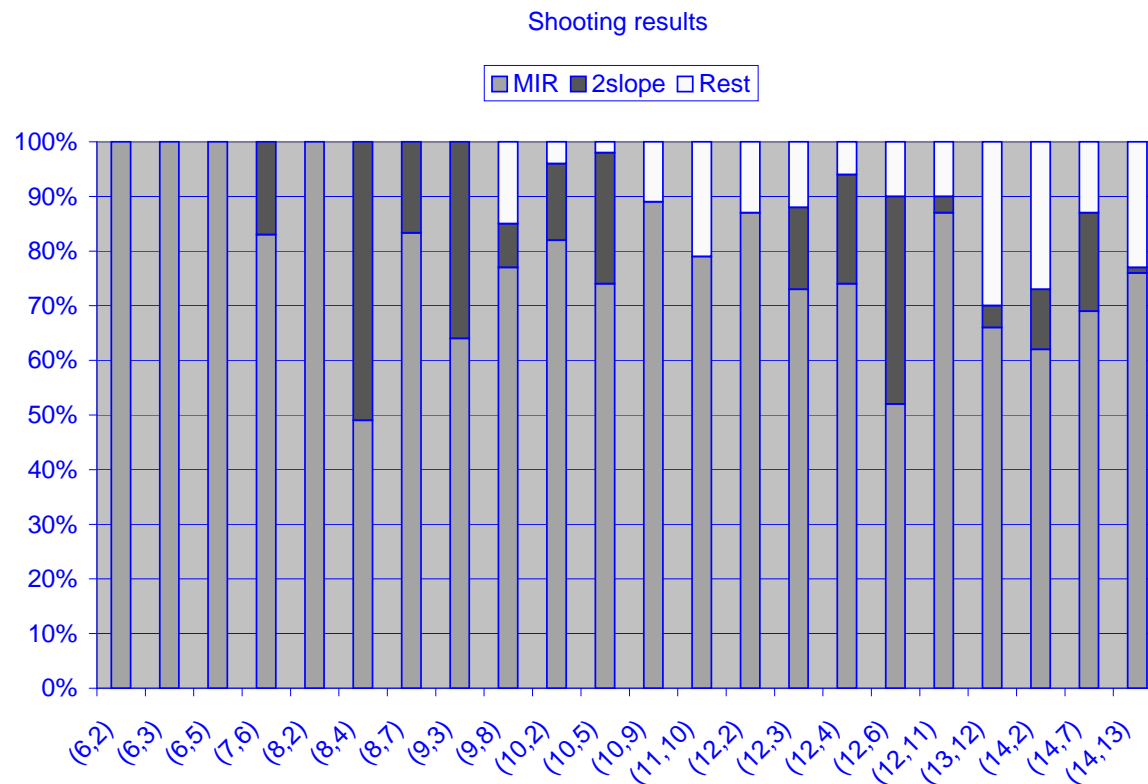
- Repeat the following experiment a number of times:
  - Choose a random direction  $d \in R_+^{n-1}$
  - Find the facet of  $P(n, r)$  first encountered along the ray  $\{\lambda d \mid \lambda \geq 0\}$
  - Record the facet.
- If a facet is hit many times, then consider it as an important facet.

- Random means  $d/\|d\|_2$  is uniformly distributed over the surface of the unit-sphere.
- Hit count is proportional to the size of the “solid angle”
- The experiment assumes that
  - Large facets are more important than the rest
  - Large facets subtend large angles at the origin.



# Shooting results of Gomory, Johnson, Evans (2003)

- For small  $n$  ( $< 30$ ), they repeat the experiment 10,000 times
- Observe that a small number of facets absorb significant number of hits



# Our Shooting Results

- For larger  $n$  ( $< 200$ ), repeat the experiment 100,000 times
- Focus on MIR-based facets

- Main observation: MIR-based facets still appear important for large  $n$ .

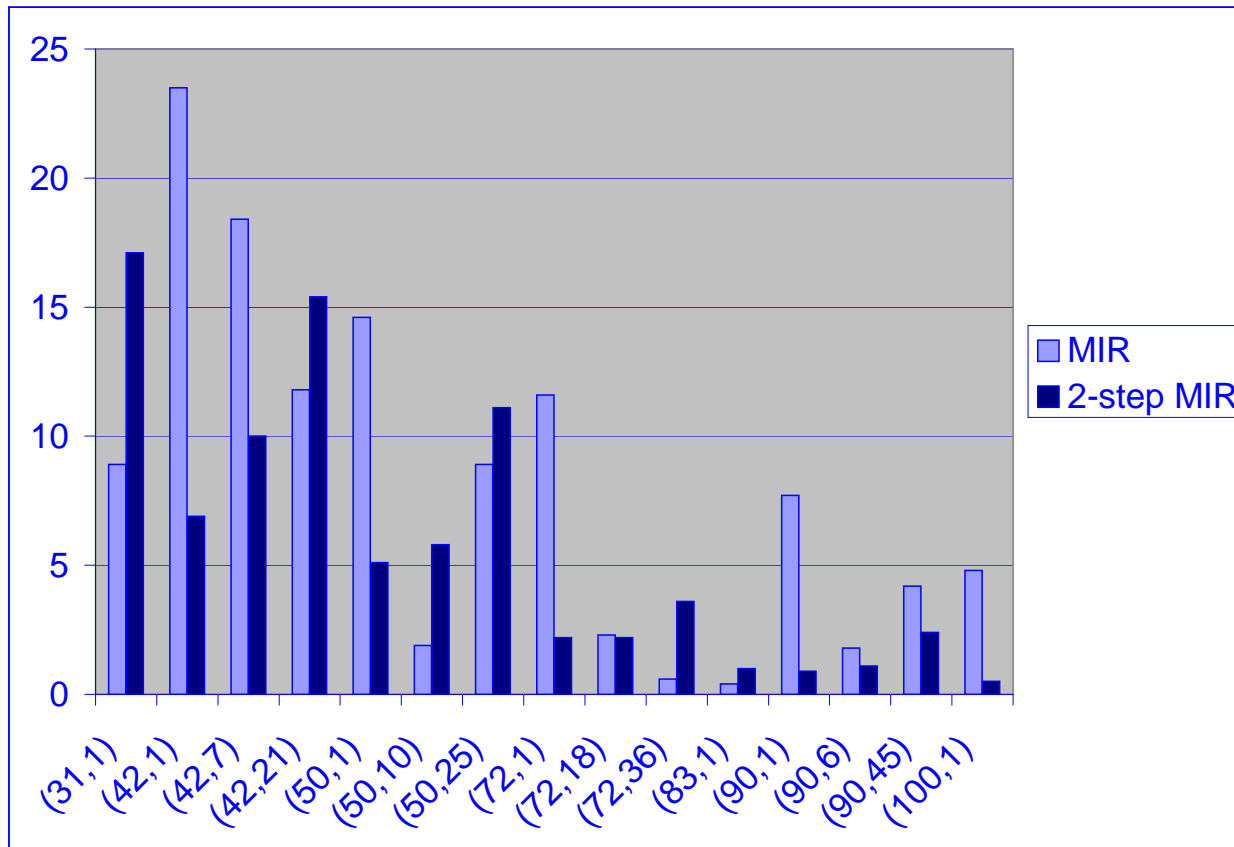
the group	total MIRs	total facets hit	% all mirs	% top 10	% top 10 MIRs
$P(31,1)$	194	29,617	26.0	8.6	8.6
$P(42,1)$	207	46,407	30.4	21.9	21.8
$P(42,7)$	216	47,918	28.4	19.3	19.2
$P(42,21)$	247	53,754	27.2	15.5	15.5
$P(50,1)$	378	65,346	19.7	13.7	13.7
$P(50,10)$	370	74,736	7.6	1.4	1.3
$P(50,25)$	465	68,202	20.0	11.2	11.2

Table 1: Shots absorbed by MIR-based facets for selected  $P(n, r)$

- but the hit rate becomes less dramatic as  $n$  increases.

## Shooting results and two-step MIR's

- Two-step MIR facets appear to be the second most important class.
- There are at most  $(n/2)$  MIR facets for  $P(n, r)$
- There are at most  $(n^2/2)$  two-step MIR facets for  $P(n, r)$



# Comparing MIR and 2-step MIR functions for $W$

- Applied to  $v_1 - v_2 + \sum_{i \in I} a_i x_i \geq b$  two-step MIR inequality

$$\frac{v_1}{\rho\tau} + \sum_{j \in I} [a_j] x_j + \frac{1}{\rho\tau} \sum_{j \in I} \min\{\rho\tau, k_j\rho + \hat{a}_j - k_j\alpha, l_j\rho\} x_j \geq [b],$$

- MIR inequality

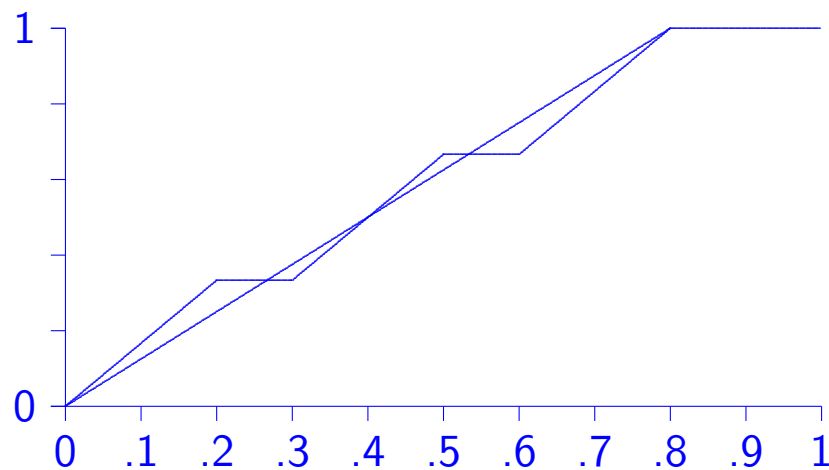
$$\frac{v_1}{\hat{b}} + \sum_{j \in I} [a_j] x_j + \frac{1}{\hat{b}} \sum_{j \in I} \min\{\hat{b}, \hat{a}_j\} x_j \geq [b],$$

where  $\hat{b} = \rho + (\tau - 1)\alpha \Rightarrow \hat{b} - \rho\tau = (\alpha - \rho)(\tau - 1)$

## Example:

$$\hat{b} = 0.8, \alpha = 0.3$$

$$\rho = 0.2, \tau = 3$$



Given

$$(\bar{v}, \bar{x}) \in W^R = \left\{ v, x \in R^{2+|I|} : v_1 - v_2 + \sum_{i \in I} a_i x_i \geq b, \quad x, v \geq 0 \right\},$$

how to choose an  $\alpha$  that gives a violated inequality?

The violation of the two-step MIR inequality generated by

$$\alpha \in \mathcal{I}^\tau = \begin{cases} (\hat{b}/\tau, 1/\tau] & \text{if } 2 \leq \tau < \lceil 1/(1 - \hat{b}) \rceil \\ (\hat{b}/\tau, \hat{b}/(\tau - 1)) & \text{if } \tau \geq \lceil 1/(1 - \hat{b}) \rceil. \end{cases} \quad (1)$$

is

$$f(\alpha) = \rho^\alpha \tau \lceil b \rceil - \bar{v}_1 - \sum_{j \in I^+} h^{a_j}(\alpha) \bar{x}_j$$

where  $\rho^\alpha = \hat{b} - \alpha(\tau - 1)$  and

$$h^a(\alpha) = \rho^\alpha \tau \lceil a \rceil + \min\{\rho^\alpha \tau, \lceil \hat{a}/\alpha \rceil \rho^\alpha + \hat{a} - \lceil \hat{a}/\alpha \rceil \alpha, \lceil \hat{a}/\alpha \rceil \rho^\alpha\}.$$

It is possible to show the following for  $\alpha \in \mathcal{I}^\tau$ :

1. The cut coefficient function  $h^a(\alpha)$  is continuous.
2. The violation function

$$f(\alpha) = \rho^\alpha \tau \lceil b \rceil - \bar{v}_1 - \sum_{j \in I^+} h^{a_j}(\alpha) \bar{x}_j$$

is continuous.

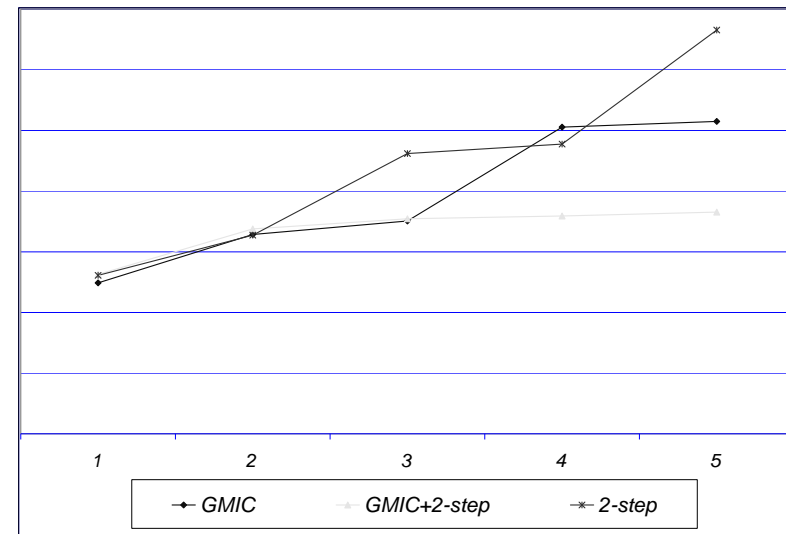
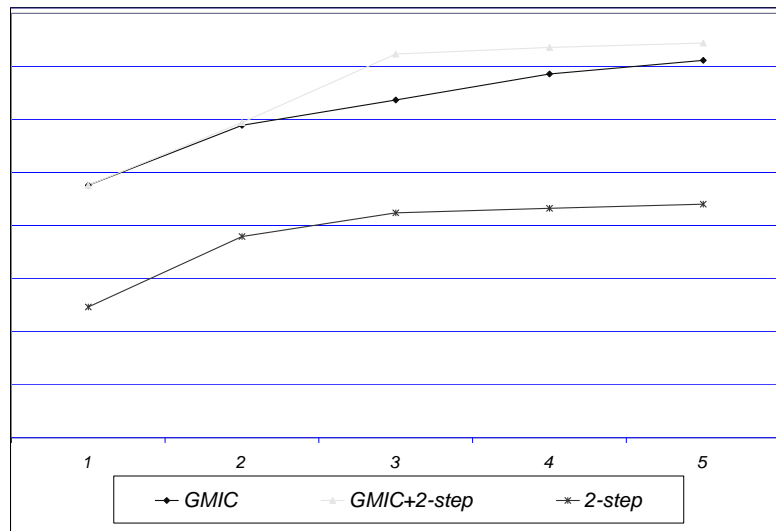
3. The violation function  $f(\alpha)$  is piecewise linear.
4. Breakpoints of the function  $f(\alpha)$  happen at  $\alpha = \hat{a}_i/t$  for some  $i \in I^+$  and  $t \in \mathbb{Z}_+$ .
5. There are at most  $|I^+|$  such  $\alpha$ 's and they can be identified by inspection.
6. The best  $\alpha \in \mathcal{I}^\tau$  can be identified in linear-time.
7. The best  $\alpha \in \bigcup_{2 \leq \tau \leq p} \mathcal{I}^\tau$  can be identified in  $O(p \cdot |I|)$  time.



- *How to obtain base inequalities?*
  - *Rows of the simplex tableau*
  - *Rows of the formulation (work-in-progress)*
- *Problem instances:*
  - *Unbounded Atamturk instances (random, dense, mixed-integer)*
  - *Bounded Atamturk instances (random, dense, mixed-integer)*
  - *Cornuejols, Li and Vandenbussche type instances (random, sparse, pure integer)*
  - *Dash, Kalagnanam, Reddy and Song steel instances (real, sparse, pure integer)*
  - *MIPLIB (work-in-progress)*

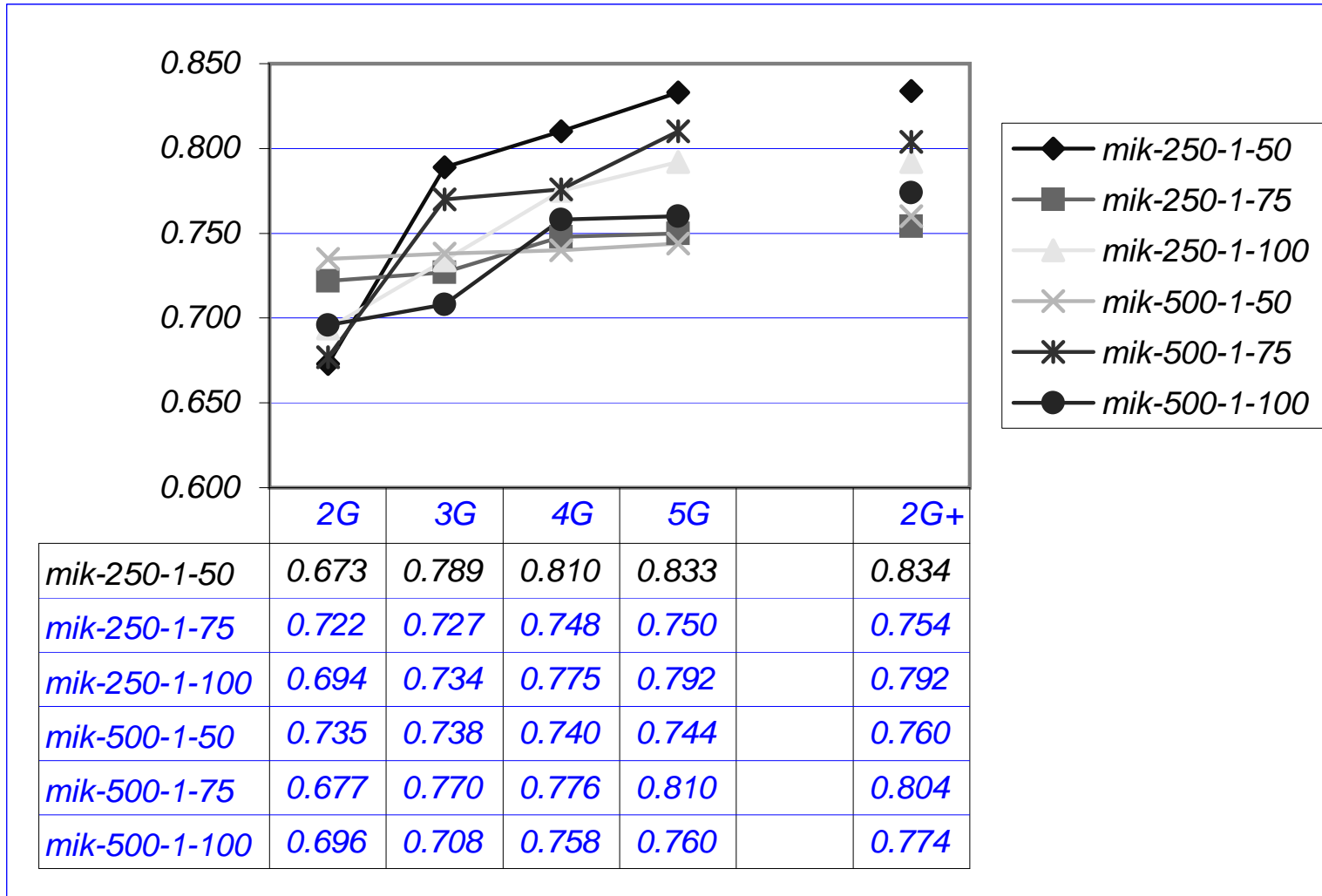
# How to measure effectiveness?

- *Two randomly generated problems with 20 rows and 100 cols:*



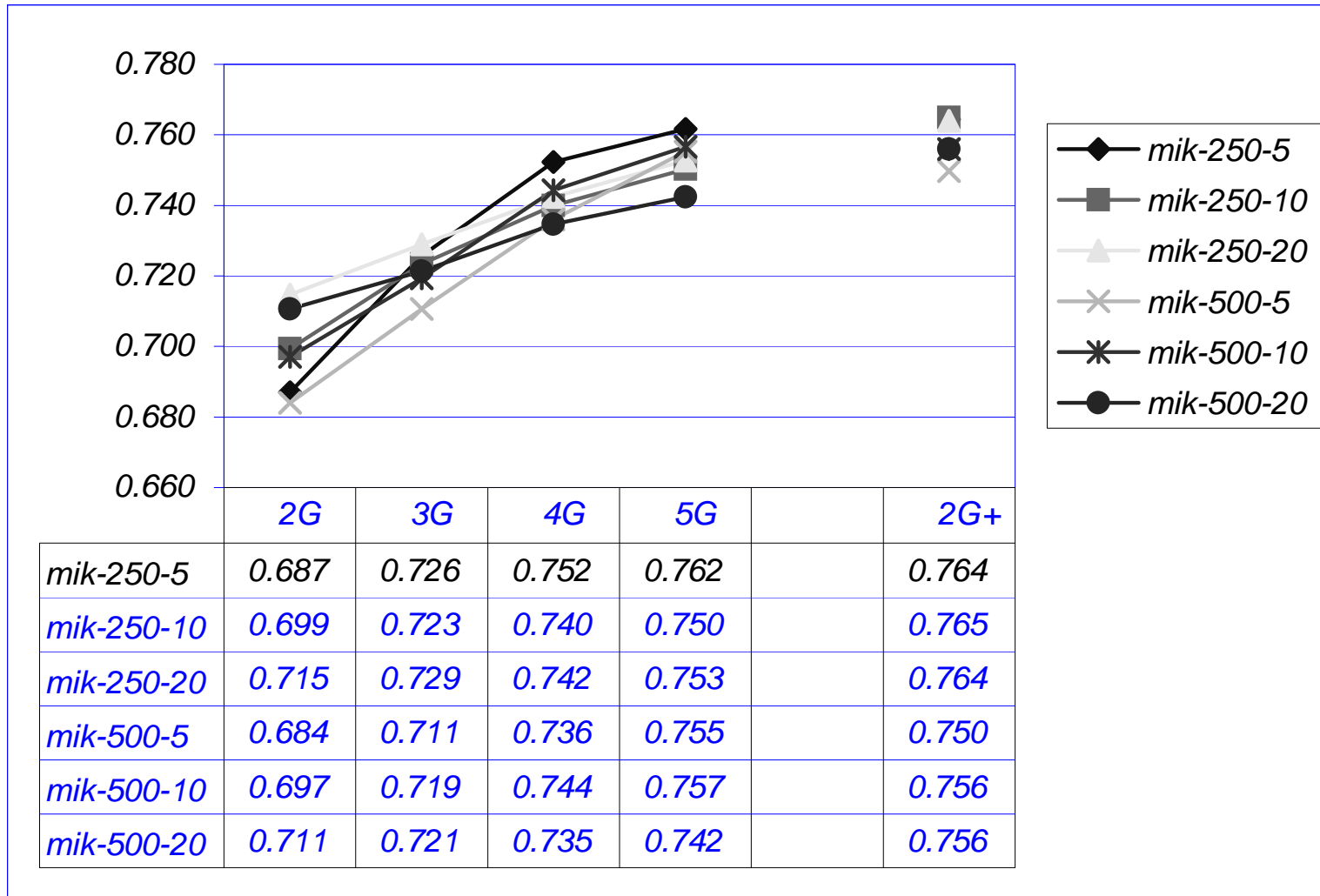
- *List cuts:*
  - *Remember the base inequalities for MIR's*
  - *Add 2-steps MIR's at the end only.*

# Unbounded Atamtürk Instances



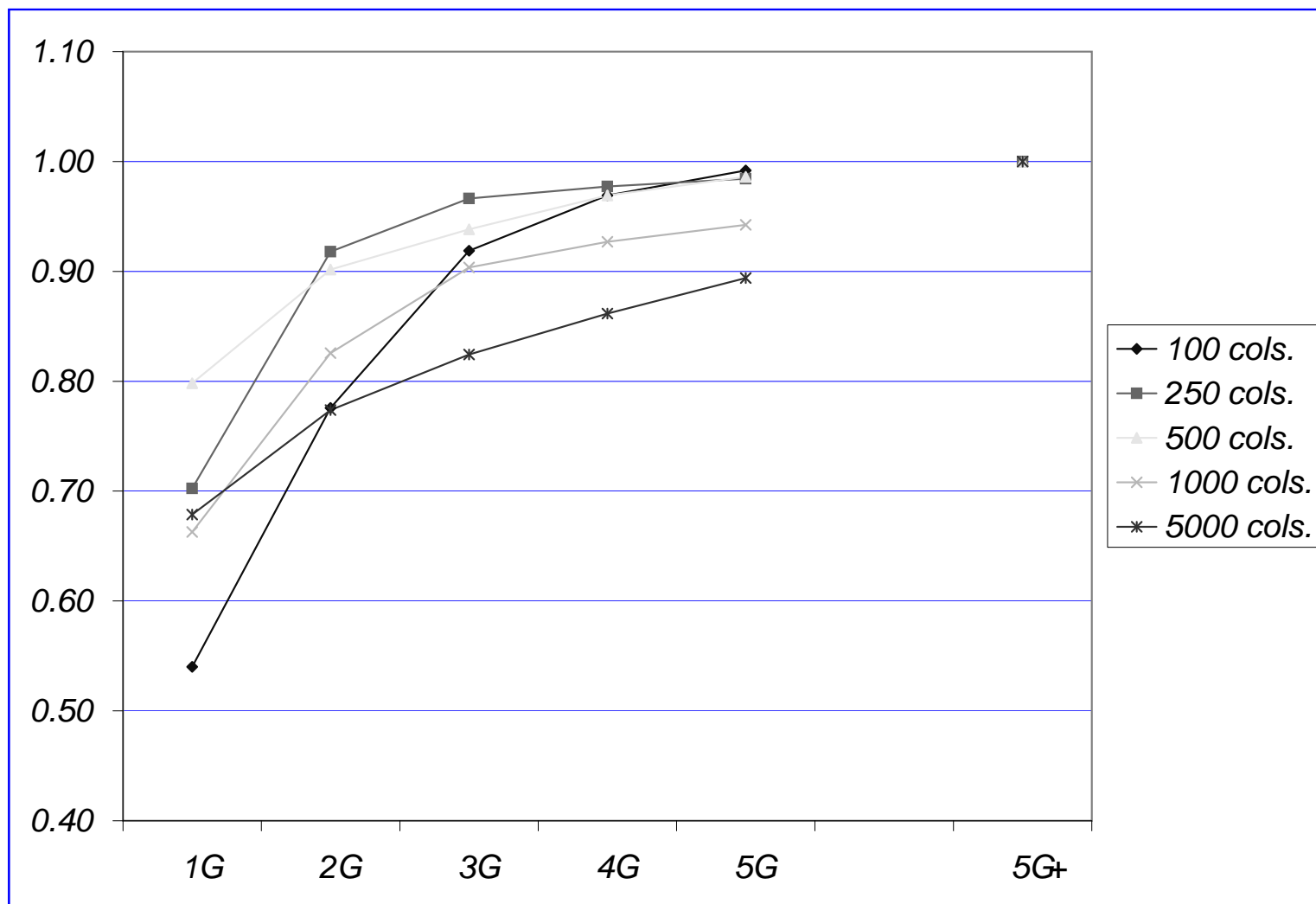
- *2G+* : 2 rounds of GMIC followed by 2-step list cuts.

## Bounded Atamtürk Instances

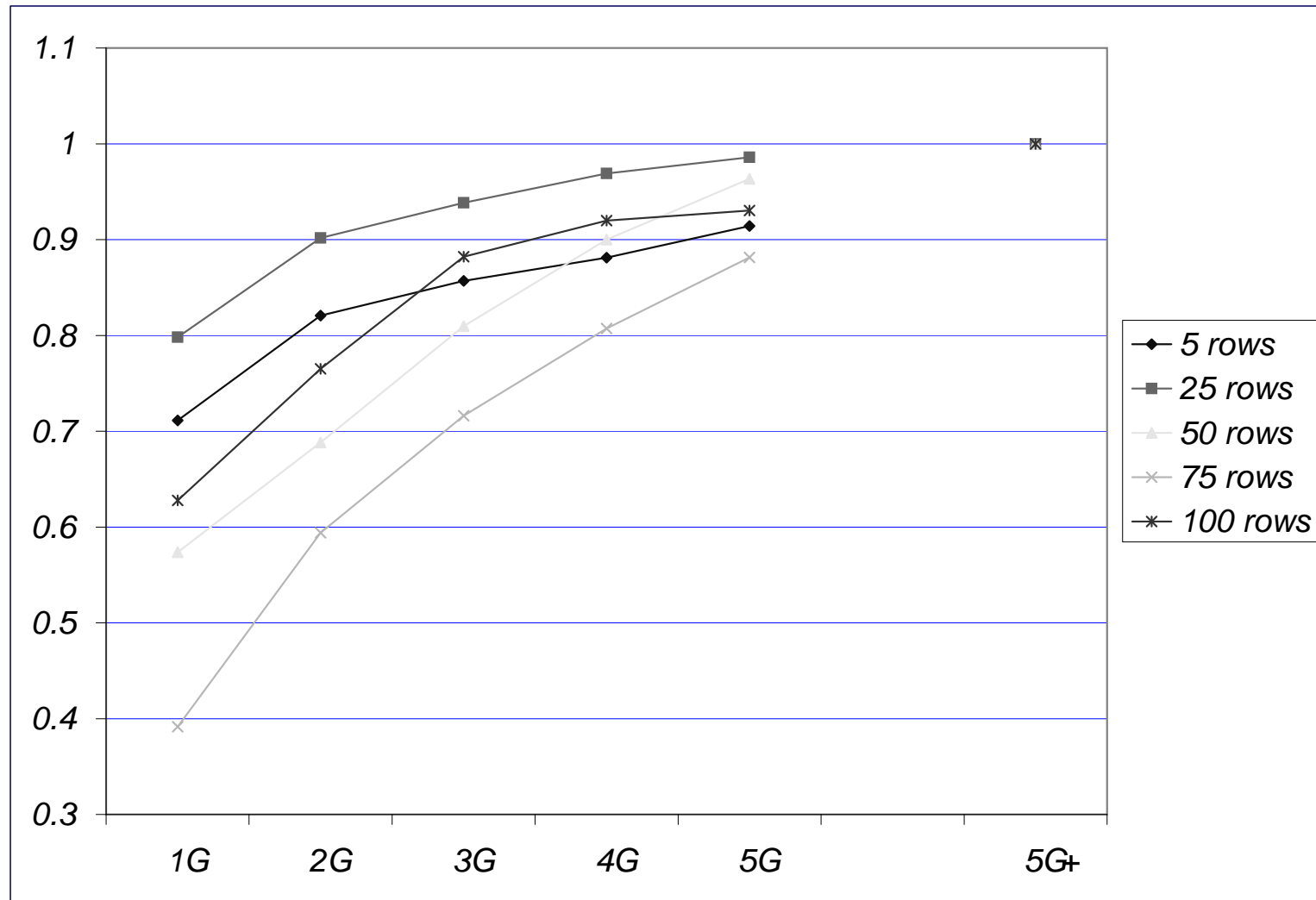


- *2G+* : 2 rounds of GMIC followed by 2-step list cuts.

## Random Instances with 25 rows

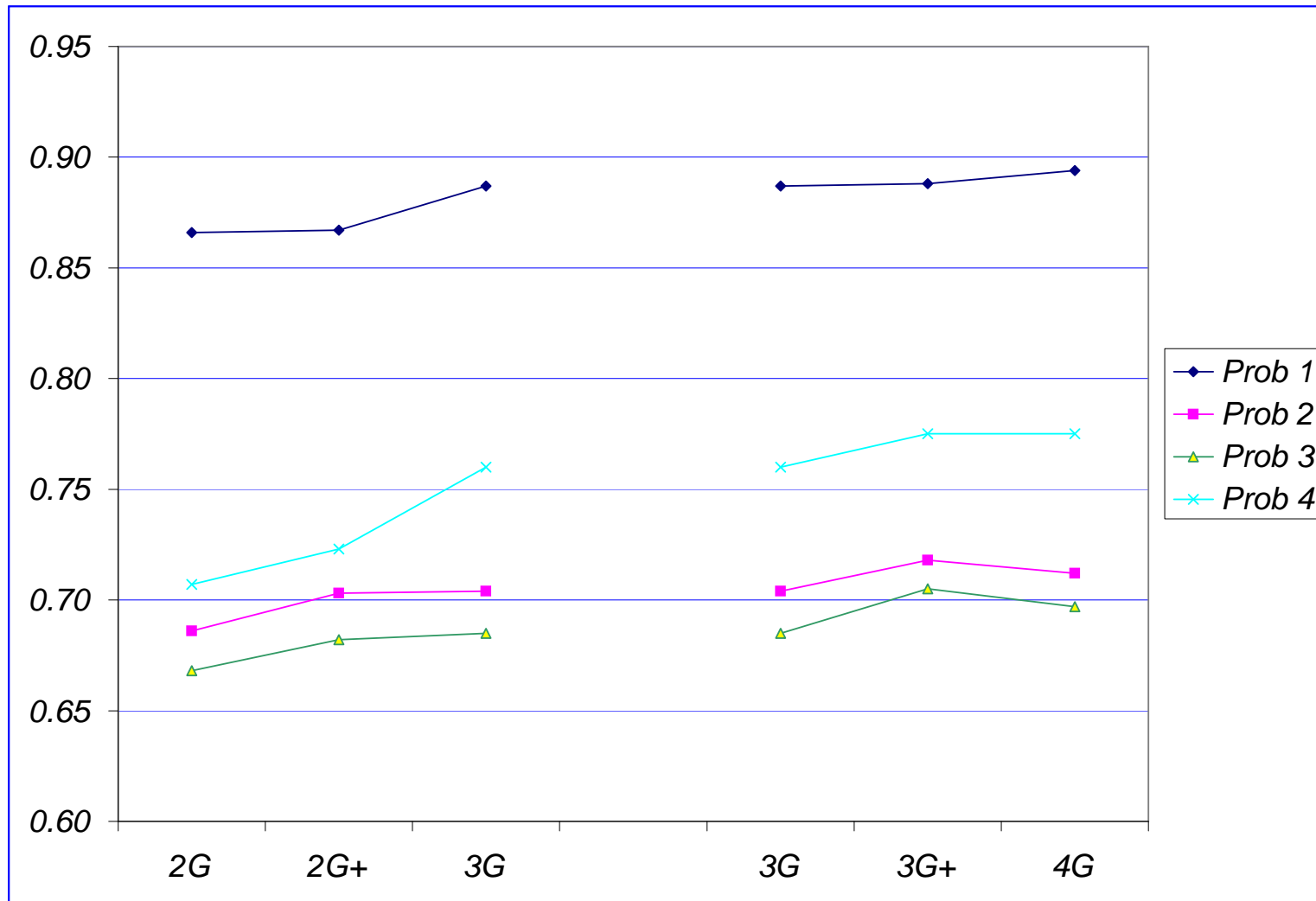


# Random Instances with 500 cols



# Steel Instances

*Cutting stock problem instances with 5000 cols and 3000 rows*



**Thank You**