

# A “finite state” version of the Kraft-McMillan theorem

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## Outline

- The Kraft-McMillan theorem
- The question of enumerative sequences of leaves
- Super-state automata
- The question of enumerative sequences of nodes

## The Kraft McMillan theorem

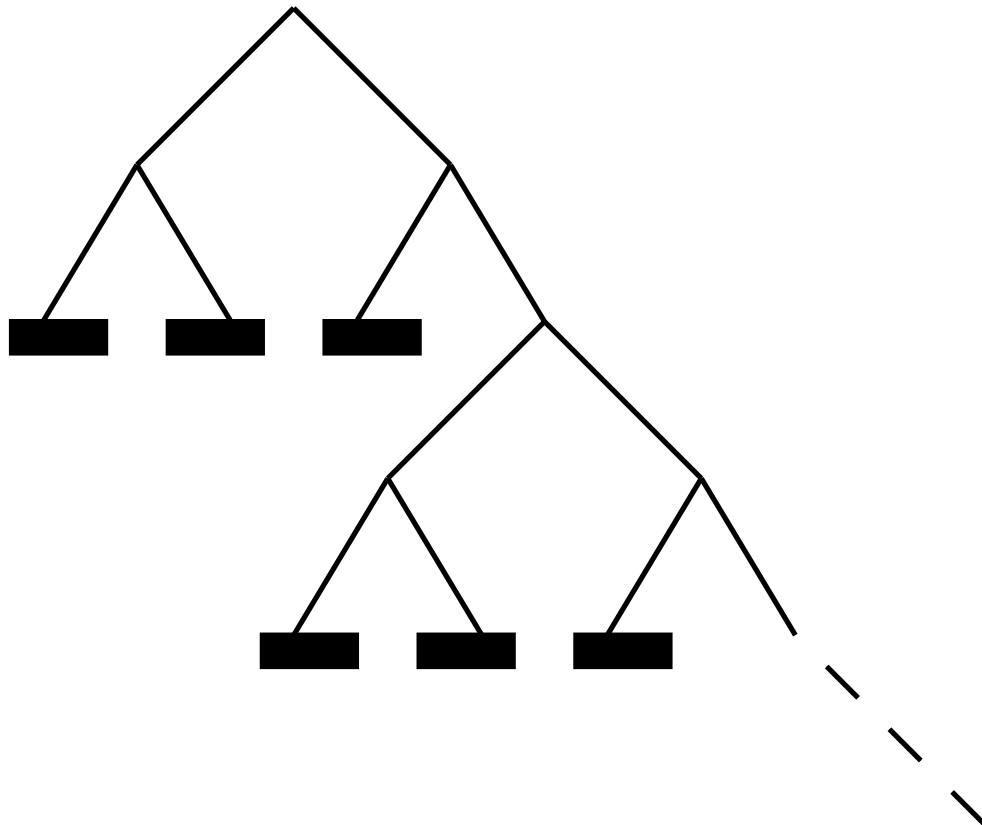
A sequence of nonnegative integers  $(s_n)_{n \geq 0}$  is the enumerative sequence of leaves of a  $k$ -ary tree (resp. a complete  $k$ -ary tree) iff

$$\sum_{n \geq 0} s_n k^{-n} \leq 1 \text{ (resp. } = 1)$$

$$s(z) = \sum_{n \geq 0} s_n z^n, \quad s\left(\frac{1}{k}\right) \leq 1$$

$s_n$  is the number of leaves at height  $n$ .

- compression coding schemes
- coding for constrained channels
- construction of prefix codes



Tree associated to  $3z^2(z^2)^* = 3z^2/(1 - z^2)$ .

## Problem

Characterize the enumerative sequences of leaves of  $k$ -ary rational trees.

**rational tree**: a tree which has a finite number of non isomorphic subtrees.

**$k$ -ary** : each node of the tree has at most  $k$  sons.

## Necessary conditions

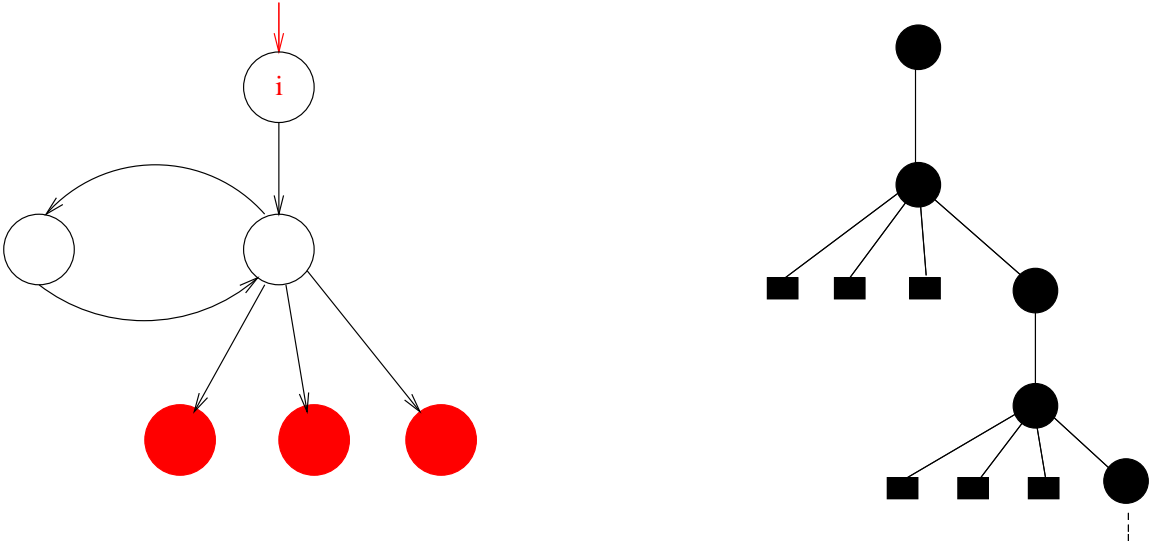
- Kraft inequality:  $\sum_{n \geq 0} s_n k^{-n} \leq 1$
- $\mathbb{N}$ -rational sequence:  $s_n$  is the number of paths of length  $n$  going from an initial state to a final state in an automaton (a finite state machine).

In general, the tree obtained with the Kraft-McMillan construction is not necessarily rational.

## Another construction

We begin with a normalized automaton that represents the sequence. We “develop” the automaton and get a rational tree. This tree is not necessarily  $k$ -ary.

$$s(z) = 3z^2(z^2)^* = 3z^2/(1 - z^2)$$



## Result

Any  $\mathbb{N}$ -rational sequence satisfying the Kraft inequality is the enumerative sequence of leaves by height of a  $k$ -ary rational tree.

## Perron-Frobenius

Let  $M$  be a nonnegative irreducible matrix. Its spectral radius  $\lambda$  is a simple eigenvalue of  $M$ . There is a positive eigenvector  $\mathbf{v}$  such that  $M\mathbf{v} = \lambda\mathbf{v}$ .

If  $\lambda < k$ , where  $k$  is a positive integer, there is an approximate eigenvector, that is a **positive** vector  $\mathbf{v}$  with  $M\mathbf{v} \leq k\mathbf{v}$ .

## Computation of an a. eigenvector

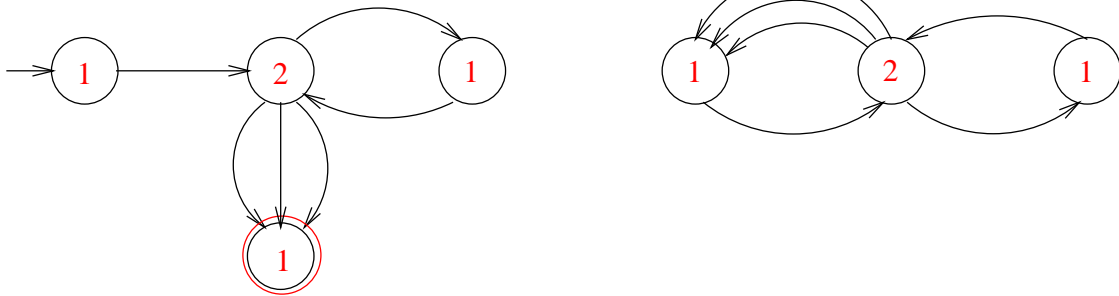
Franaszek's algorithm

```
begin  
   $\mathbf{v} := (1, 1, \dots, 1)^t$ ;  
  repeat  
    begin  
       $\mathbf{w} := \mathbf{v}$ ;  
       $\mathbf{v} := \max \left\{ \mathbf{w}, \left\lceil \frac{1}{k} M \mathbf{w} \right\rceil \right\}$ ;  
    end  
  until  $\mathbf{v} = \mathbf{w}$ ;  
end
```

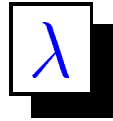
## From the series to the (a.) eigenvector

Let  $s$  be the sequence recognized by a normalized automaton  $(G, i, F)$ . Let  $\bar{G}$  be the strongly connected graph obtained from  $G$  by identification of initial and final states.

$$S(z) = 3z^2(z^2)^* = \frac{3z^2}{1-z^2}$$



$$M = \begin{pmatrix} 0 & 1 & 0 \\ 3 & 0 & 1 \\ 0 & 1 & 0 \end{pmatrix} \quad \mathbf{v} = \begin{pmatrix} 1 \\ 2 \\ 1 \end{pmatrix} \quad M\mathbf{v} = 2\mathbf{v}$$



- does not depend on the representation  $G$ . ( $\frac{1}{\lambda}$  is the minimal modulus of the poles of the series  $\frac{1}{1-s(z)}$ ).

$$\sum_{n \geq 0} s_n k^{-n} < 1 \text{ si } \lambda < k$$

$$\sum_{n \geq 0} s_n k^{-n} = 1 \text{ si } \lambda = k$$

We take a  $k$ -approximate eigenvector  $\mathbf{v}$ :  
 $M\mathbf{v} \leq k\mathbf{v}$ .

## Super-state automata

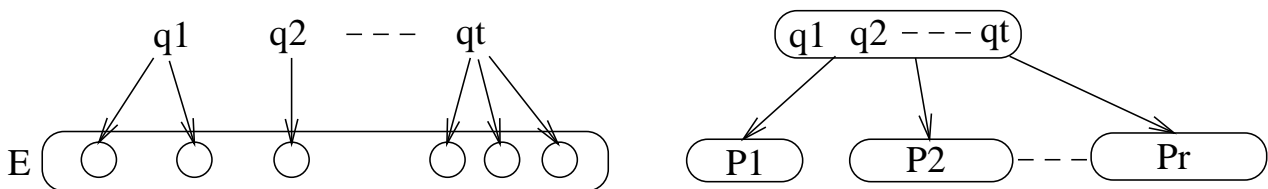
$$\mathcal{A} = (Q, F) \longrightarrow \mathcal{B} = (Q', F')$$

+ valuation  $v$  of the states + integer  $m$

$$Q' = \{(q_1, q_2, \dots, q_t) \mid q_i \in Q, 1 \leq t \leq m\}$$

set of unordered  $t$ -uples.

$$v((q_1, q_2, \dots, q_t)) = \sum_{j=1}^t v(q_j)$$



$P_1, P_2, \dots, P_r$  is a partition of  $E$ .

The partition is such that all parts but one have a valuation dividable by  $m$ .

## Weight lemma or pigeon-hole

[Brian Marcus [2]]

*Let  $v_1, v_2, \dots, v_m$  be integers. There is a subset  $S$  of  $\{1, 2, \dots, m\}$  such that  $\sum_{q \in S} v_q$  is dividable by  $m$ .*

Paul Erdős attributes this nice application of the pigeon-hole principle to Marta Sved and Andrew Vázsonyi in [1] p. 125.

## References

- [1] Martin Aigner and Günter M. Ziegler. *Proofs from The Book*, Springer, 1998.
- [2] Brian Marcus. Factors and extensions of full shifts *Monats. Math.*, 88:239-247, 1979

proof:

$$m \left\{ \begin{array}{l} v_1 \\ v_1 + v_2 \\ v_1 + v_2 + v_3 \\ \vdots \\ v_1 + v_2 + \dots + v_m \end{array} \right.$$

either there is a  $p$  such that, modulo  $m$ ,

$$v_1 + v_2 + \dots + v_p \equiv 0$$

or there are  $p, r$  such that, modulo  $m$ ,

$$v_1 + v_2 + \dots + v_p \equiv v_1 + v_2 + \dots + v_r$$

## Proof

$s \longrightarrow G \longrightarrow \bar{G} \longrightarrow \mathbf{v} : M\mathbf{v} \leq k\mathbf{v} \ (\lambda \leq k)$

One takes  $\mathbf{v}$  as **valuation** and  $m = v(i)$ .

Let  $u$  be a super-state. One defines:

$$w(u) = \lceil v(u)/m \rceil$$

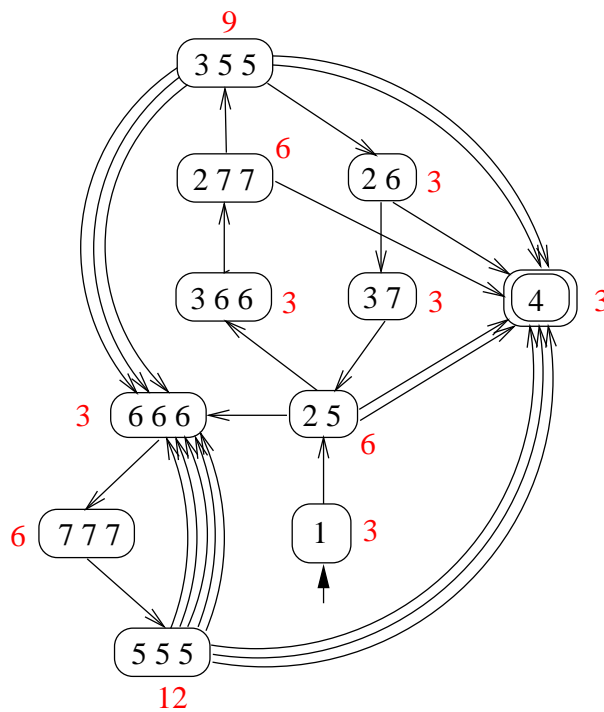
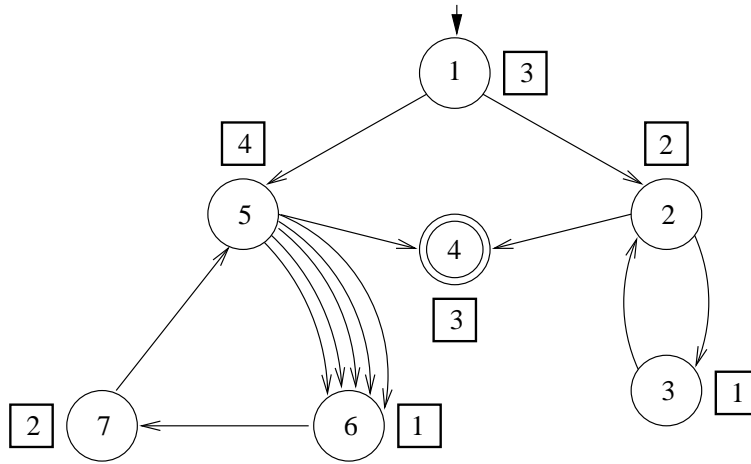
If  $u$  has  $t$  outgoing edges that arrive in  $u_1, \dots, u_t$ ,

$$\sum_{j=1}^t v(u_j) \leq kv(u)$$

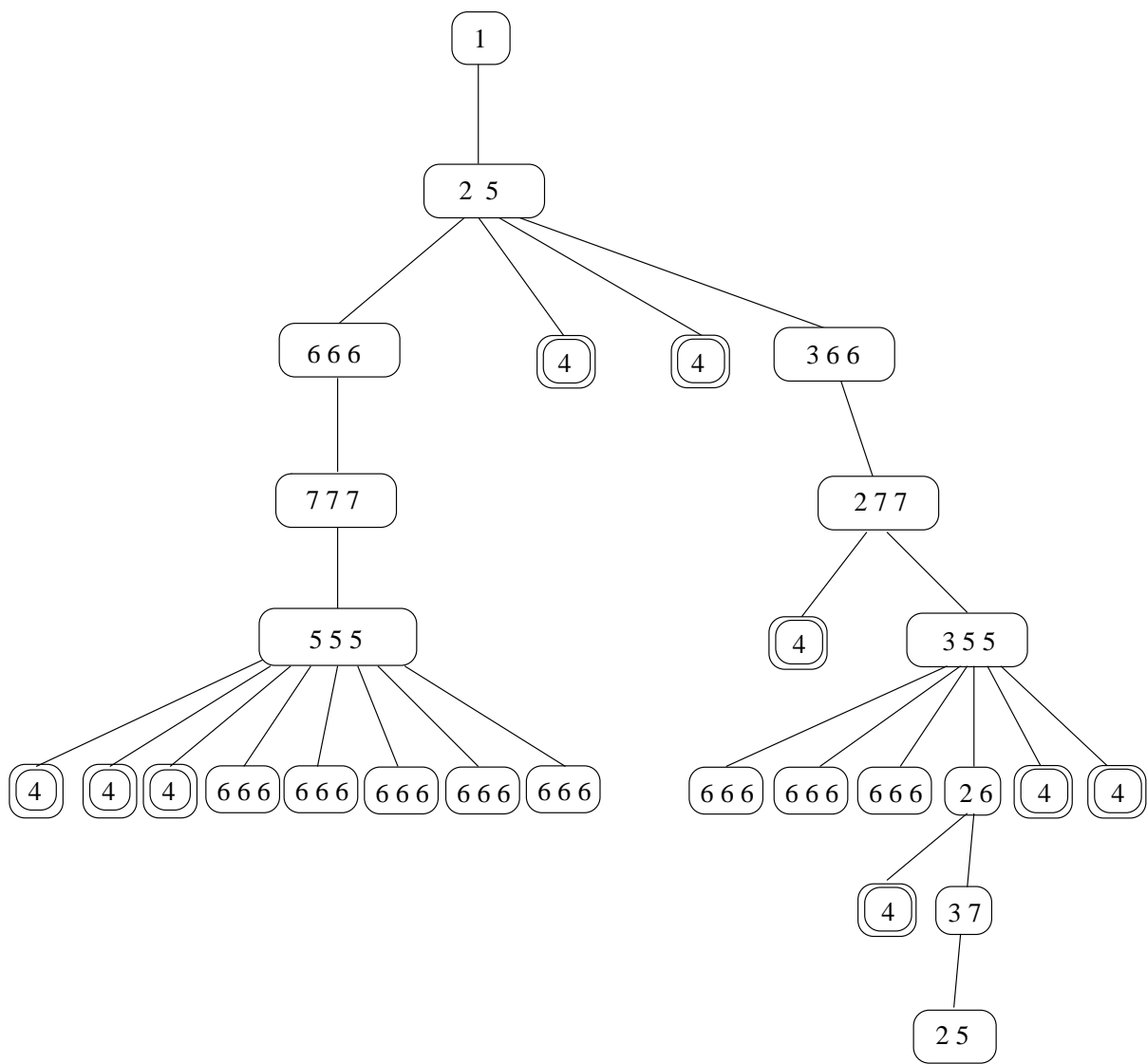
$$\sum_{j=1}^{t-1} v(u_j)/m + \lceil v(u_t)/m \rceil \leq k \lceil v(u)/m \rceil$$

$$\sum_{j=1}^t w(u_j) \leq kw(u) \text{ et } w(i) = 1$$

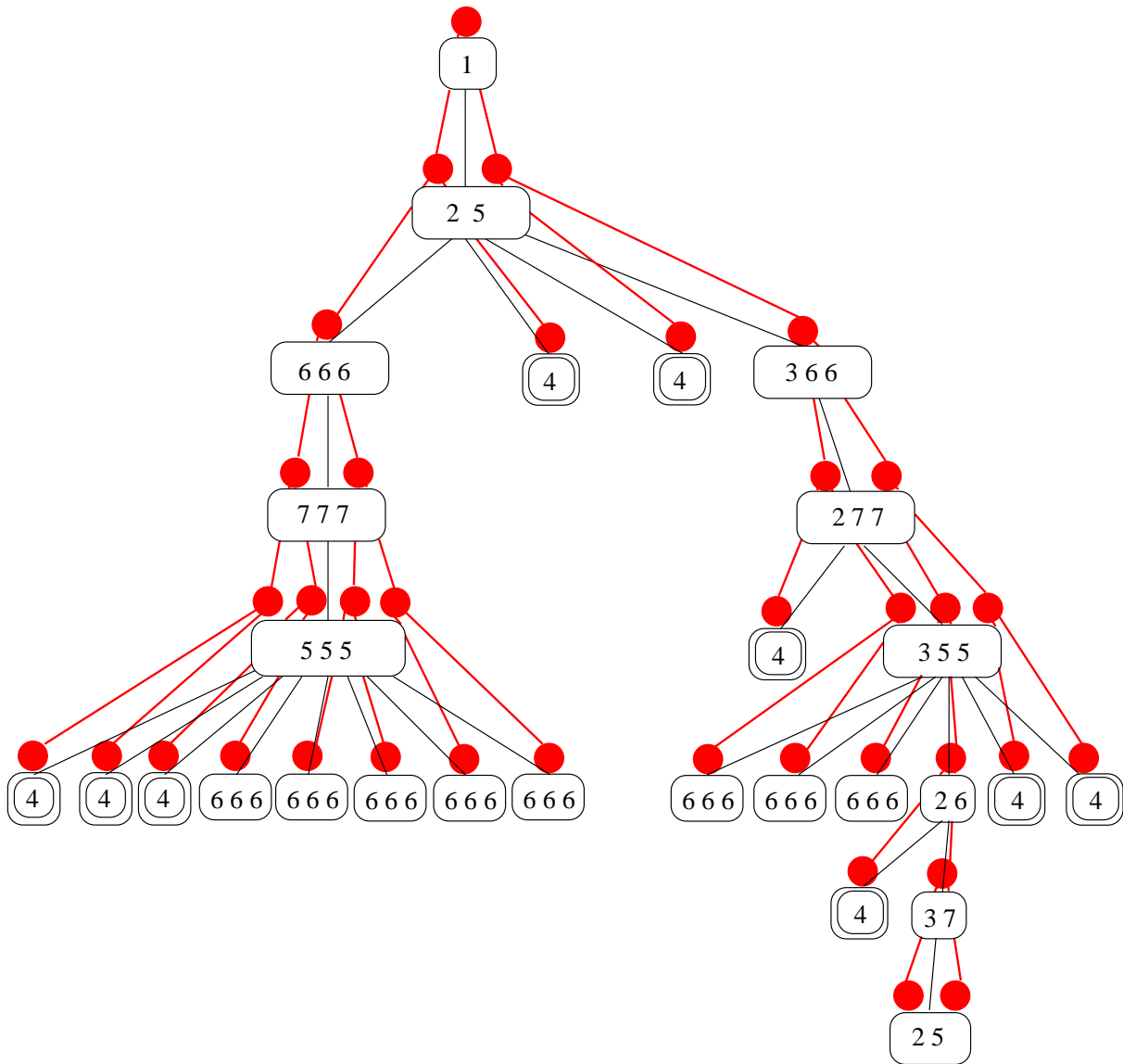
$$s(z) = \frac{z^2}{1-z^2} + \frac{z^2}{1-5z^2}$$



Super-state automaton with the **valuations**.



The developed automaton.



The number of small red balls above a super-state  $u$  is  $w(u) = \lceil \frac{v(u)}{m} \rceil$  (here  $m = 3$ ).

## The sequences of nodes

An  $\mathbb{N}$ -rational sequence  $(t_n)_{n \geq 0}$  is the enumerative sequence of nodes by height of a  $k$ -ary rational tree iff:

- $t_0 = 1$
- $\forall n \geq 1, t_{n+1} \leq kt_n$
- the convergence radius of the series  $t$  is  $> \frac{1}{k}$
- the series  $s(z) = t(z)(kz - 1) + 1$  is  $\mathbb{N}$ -rational.

There are  $\mathbb{N}$ -rational series that satisfy the first three conditions but not the last one.

# References

- [1] F. Bassino, M.-P. Béal, and D. Perrin. Super-state automate and rational trees. In *LATIN'98*, LNCS 1380, pages 42–52. Springer, 1998.
- [2] F. Bassino, M.-P. Béal, and D. Perrin. Enumerative sequences of leaves and nodes in rational trees. *Theoret. Comput. Sci.*, 1997. (to appear).
- [3] F. Bassino, M.-P. Béal, and D. Perrin. Enumerative sequences of leaves in rational trees. In *ICALP'97*, LNCS 1256, pages 76–86. Springer, 1997.
- [4] D. Perrin. Arbres et séries rationnelles. *C.R.A.S. Paris, Série I*, 309:713–716, 1989.
- [5] D. Perrin. A conjecture on rational sequences. In R. Capocelli, editor, *Sequences*, pages 267–274. Springer-Verlag, 1990.