

A DETERMINISTIC · CONTROL · BASED

APPROACH TO

MOTION BY CURVATURE

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JOINT WORK WITH SYLVIA SERPACU

PREPRINT: ON MY WEB PAGE

CRUCIAL INGREDIENT : G. BARLES + F. DA LIO

PRIOR, RELATED WORK: CATTE', DIBOS, KOEFLER
GUICHARD
"MORPHOLOGICAL EVOLUTIONS"

INTERFACE MOTION

- MATERIALS SCIENCE: PHASE BOUNDARIES, GRAIN BOUNDARIES, ETC
- IMAGE PROCESSING: EDGES; DENOISING VIA NONLINEAR DIFFUSION
- HIGHER CODIMENSION ALSO INTERESTING: E.G. CURVES IN \mathbb{R}^3

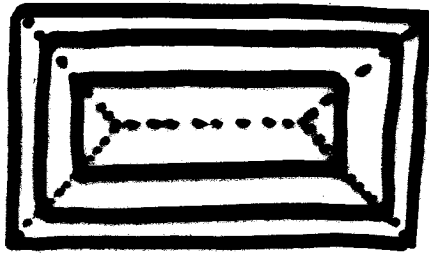
LEVEL SET METHOD

- FLEXIBLE NUMERICAL TOOL
- POWERFUL FOR ANALYSIS TOO

THIS TALK

- FRESH UNDERSTANDING VIA OPT'L CONTROL

REVIEW FAMILIAR FACTS ABOUT $v = 1$



$\Omega = \text{RECTANGLE}$

ARRIVAL TIME

$u(x)$ = TIME OF ARRIVAL AT x
SOLVES $|v| = 1$, $u = 0$ AT $\partial\Omega$

LEVEL SET METHOD

IF $v_x / |v| = -1$ THEN EACH
LEVEL SET OF v HAS VELOCITY ± 1

EQUIVALENCE

$$v(x, t) = u(x) - t$$

1st ORDER HTE

METHOD OF CHARACTERISTICS;
NOTE SHOCKS

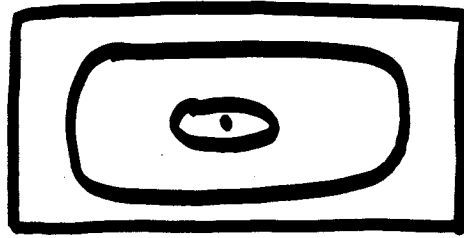
KINETIC

MODELS DECAY (OR GROWTH)

OPTIMAL CONTROL

$u(x) = \min_{\text{VELOCITY} \leq 1} \{ \text{ARRIVAL TIME TO } \partial\Omega \}$
STARTING AT x

REVIEW FAMILIAR FACTS ABOUT $v = -K$



CONVEX Ω

ARRIVAL TIME

$u(x)$ = TIME OF ARRIVAL AT x
SOLVES $\operatorname{div}(\nabla u / |\nabla u|) = -1 / |\nabla u|$

LEVEL SET METHOD

IF $v_t / |\nabla v| = \operatorname{div}(\nabla v / |\nabla v|)$ THEN
EACH LEVEL SET OF v HAS VELOCITY $= -K$

EQUIVALENCE

$$v(x, t) = u(x) - t$$

2nd ORDER

NO CHARACTERISTICS

THERMODYNAMIC

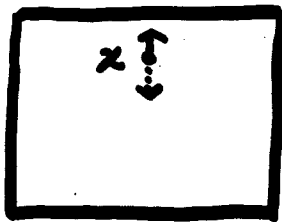
STEEPEST DESCENT
FOR PERIMETER

OPTIMAL CONTROL

2nd ORDER \Leftrightarrow STOCHASTIC CONTROL

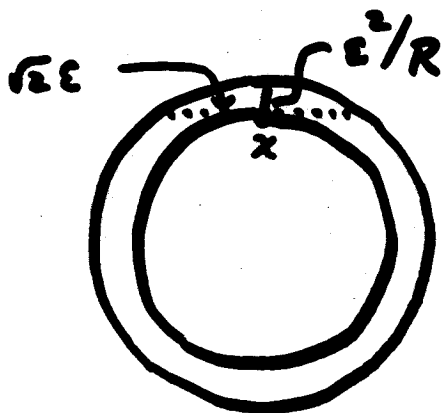
THE GAME

PAUL WANTS TO EXIT; CAROL WANTS TO STOP HIM; STEP SIZE IS $\epsilon > 0$.

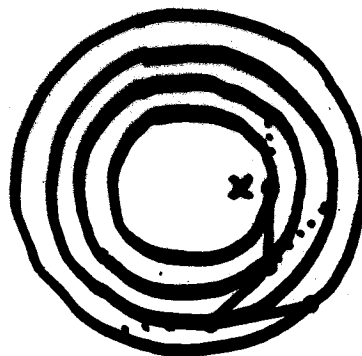


- PAUL CHOOSES DIR'N $|v| = 1$
- CAROL MAY REVERSE IT $b = \pm 1$
- PAUL GOES $x \rightarrow x + \sqrt{\epsilon} b v$

CAN PAUL EXIT? YES!

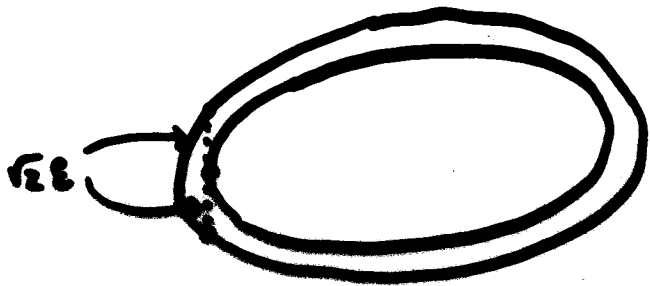


RADIUS R



- SET FROM WHICH PAUL CAN EXIT IN I STEPS HAS NORMAL VELOCITY $1/R_I$ (AFTER NORMALIZATION)
- CAROL CAN'T STOP PAUL, BUT SHE DOES SLOW HIM DOWN

GEOMETRIC INTERPRETATION



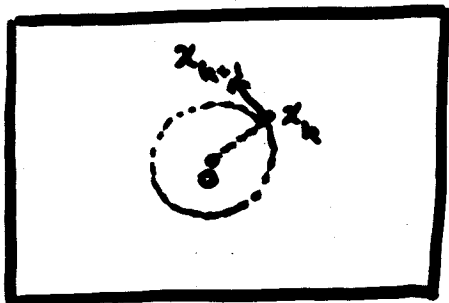
SET FROM WHICH PAUL
NEEDS TWO STEPS TO EXIT
IS LOCUS OF MIDPOINTS
OF SECANTS

NONCONVEX CASE



SAME AS ABOVE, BUT HE
CANNOT EXIT AT
CONCAVE PART OF $\partial\Omega$

HE CAN EXIT IN $\mathcal{O}(\epsilon^{-2})$ STEPS



TRIAL STRATEGY: USE
 $v \perp x$. THEN
 $|x_{k+1}|^2 = x_k^2 + 2\epsilon^2$

GOING BEYOND PICTURES

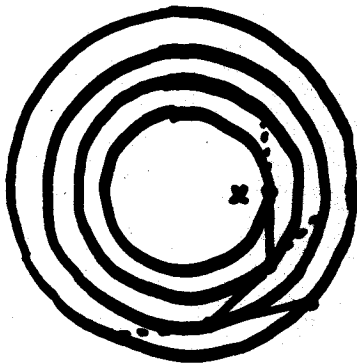
$$u_{\varepsilon}(x) = \varepsilon^2 \cdot \left[\begin{array}{l} \# \text{ STEPS PAUL NEEDS TO} \\ \text{EXIT, STARTING AT } x \end{array} \right]$$

SIMPLEST THEOREM: FOR CONVEX PLANE DOMAINS

$$\lim_{\varepsilon \rightarrow 0} u_{\varepsilon}(x) = \text{ARRIVAL TIME OF } \partial\Omega \\ \text{UNDER MOTION BY CURVATURE}$$

MAIN TOOL: DYNAMIC PROGRAMMING PRINCIPLE

$$u_{\varepsilon}(x) = \min_{|v|=1} \max_{b=\pm 1} \left\{ u_{\varepsilon}(x + \sqrt{2}\varepsilon bv) + \varepsilon^2 \right\}$$



"LEVEL-SET PDE" IS THE HJB EQU TO OPT'L CONTROL PROBLEM. FORMALY:

$$u(x) = \min_{\varepsilon} \max_{|v|=1} \max_{b=\pm 1} u_{\varepsilon}(x + \sqrt{2}\varepsilon bv) + \varepsilon^2$$

$$\Rightarrow u(x) \approx \min_{|v|=1} \max_{b=\pm 1} u(x) + \sqrt{2}\varepsilon bv \cdot \nabla u + \varepsilon^2 \langle D^2 u \cdot v, v \rangle + \varepsilon^2$$

$\max_{b=\pm 1} \varepsilon bv \cdot \nabla u = \varepsilon v \cdot \nabla u $	DOMINATES, UNLESS = 0
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So $\cancel{u(x)} \approx \min_{v \perp \nabla u} \cancel{u(x)} + \varepsilon^2 \langle D^2 u \cdot v, v \rangle + \varepsilon^2$

$$\langle D^2 u \cdot \frac{\nabla u^{\perp}}{|\nabla u|}, \frac{\nabla u^{\perp}}{|\nabla u|} \rangle + 1 = 0$$

$$\Leftrightarrow \Delta u - \langle D^2 u \cdot \frac{\nabla u}{|\nabla u|}, \frac{\nabla u}{|\nabla u|} \rangle + 1 = 0 \quad \text{in } \mathbb{R}^2$$

$$\Leftrightarrow |\nabla u| \operatorname{div} \left(\frac{\nabla u}{|\nabla u|} \right) + 1 = 0$$

PDE IS 2 nd ORDER BECAUSE 1 st ORDER TAYLOR EXPANSION WAS NOT SUFFICIENT
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EXTENSIONS

1) SAME GAME IN \mathbb{R}^3 ?

BOUNDARY MOVES WITH NORMAL VELOCITY
= LARGEST PRINCIPAL CURVATURE

2) CAN WE GET $v = \text{MEAN CURVATURE}$ IN \mathbb{R}^3 ?

YES, WITH MODIFIED GAME

- PAUL CHOOSES TWO ORTHOG. DIRNS
 $|v| = |w| = 1, v \perp w$
- CAROL MAY REVERSE EITHER (OR BOTH)
 $b = \pm 1, \beta = \pm 1$
- PAUL GOES $x \rightarrow x + \epsilon \beta b v + \epsilon \beta \alpha w$

3) WHAT IF Ω IS NOT CONVEX?



$\lim_{\epsilon \rightarrow 0} u_\epsilon(x) = \text{ARRIVAL TIME OF}$
 $\partial\Omega$ UNDER FLOW
BY VELOCITY = K_+

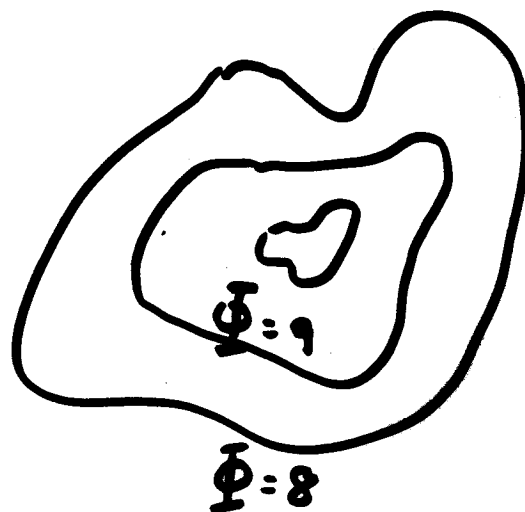
4) THE ANALOGOUS FINAL-TIME GAME?

SAME RULE BUT PAUL HAS
DIFFERENT GOAL:

$$V_{\epsilon}(x, t) = \min_{\epsilon} \Phi(y_{\epsilon}(T))$$

STARTING
AT x AT
TIME t

WHERE $y_{\epsilon}(t)$ = PAUL'S PATH
 Φ = FIXED OBJECTIVE
 T : "MATURITY TIME"



THEN

$$v_t + \langle D^2 v, \frac{\nabla v}{|\nabla v|}, \frac{\nabla v}{|\nabla v|} \rangle = 0 \quad t < T$$

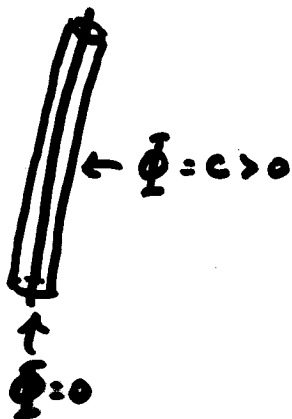
$$v = \bar{\Phi} \quad t = T$$

EACH LEVEL SET OF v MOVES
BY CURVATURE (BACKWARD IN TIME)

5) MOTION BY CURVATURE OF A CURVE IN \mathbb{R}^3

USE TIME-DEPENDENT VERSION, WITH OBJECTIVE

$$\Phi(y) = \text{dist}(y, \text{CURVE})$$



$$v_t + \min_{\substack{\xi \perp \nabla v \\ |\xi| = 1}} \langle D^2 v \cdot \xi, \xi \rangle = 0 \quad t < T$$

$$v = \bar{\Phi} \quad t = T$$

EACH LEVEL SET OF v HAS VELOCITY: SMALLER PRW. CURVATURE

AMBROSIO-SONER: ZERO-LEVEL SET "MOVES WITH VELOCITY: CURVATURE"

RIGOROUS ANALYSIS: TWO ALTERNATIVES

VISCOSITY SOLUTIONS

OR

VERIFICATION ARGUMENT

- MORE GENERAL
- CONVERGENCE NO RATE
- GUARANTEED TO SUCCEED

- CLASSICAL + ELEMENTARY
- GETS CONV. RATE
- NEEDS SMOOTHNESS

FOCUS HERE ON ONE RESULT, PROVED
VIA VERIFICATION ARGUMENT

THEOREM: FOR SMOOTHLY BOUNDED CONVEX
PLANE DOMAINS,

$$u(x) - C\varepsilon \leq u_\varepsilon(x) \leq u(x) + C\varepsilon$$

WHERE

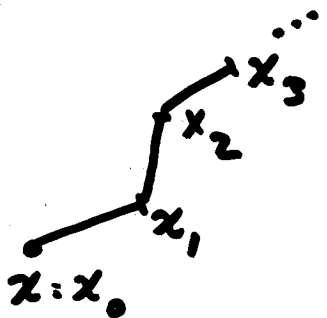
u_ε = PAUL'S SCALED ARRIVAL TIME

u : CURVATURE FLOW ARRIVAL TIME

NOTE: • UPPER BOUND REQUIRES ONE GOOD STRATEGY
• LOWER BOUND MUST CONSIDER OPT'L STRATEGY
• PROOF RESEMBLES 1st ORDER THEORY

SKETCH OF UPPER BOUND $u_\epsilon(x) \leq u(x) + C\epsilon$

TRIAL STRATEGY: ALWAYS CHOOSE $v = \pm \frac{\nabla u}{|\nabla u|}$



$$x_{k+1} = x_k + \sqrt{2} \epsilon b_k v_k$$

UNTIL EXIT AT STEP K .

$$u_\epsilon(x_k) = \min_{|v|=1} \max_{b=\pm 1} u_\epsilon(x_k + \sqrt{2} \epsilon b v) + \epsilon^2$$

$$\leq \max_{b=\pm 1} u_\epsilon(x_k + \sqrt{2} \epsilon b v_k) + \epsilon^2 \quad v_k = \frac{\nabla u}{|\nabla u|}$$

$$= u_\epsilon(x_k + \sqrt{2} \epsilon b_k v_k) + \epsilon^2$$

\Rightarrow
 $u_\epsilon(x_0) \leq K\epsilon^2$
WHERE EXIT OCCURS AT STEP K .

NOW CONSIDER $u(x)$: SINCE

$$x_{k+1} = x_k + \sqrt{2} \varepsilon b_k v_k \quad v_k = \frac{\nabla u}{|\nabla u|}(x_k)$$

WE HAVE

$$u(x_{k+1}) = u(x_k) + \cancel{\sqrt{2} \varepsilon b_k v_k \cdot \nabla u} + \underbrace{\varepsilon^2 \langle D^2 u \cdot v_k, v_k \rangle}_{-1} + \mathcal{O}(\varepsilon^3)$$

$$u(x_{k+1}) = u(x_k) - \varepsilon^2 + \mathcal{O}(\varepsilon^3)$$

\Rightarrow

$u(x_0) = K \varepsilon^2 + \mathcal{O}(\varepsilon)$	WHERE EXIT IS AT STEP K
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THUS

$$u_\varepsilon(x_0) \leq K \varepsilon^2 \leq u(x_0) + C \varepsilon \quad \text{QED}$$

NOTE: ARGUMENT NEEDS THAT

$u(x)$: ARRIVAL TIME OF $\partial\Omega$
UNDER CURVATURE FLOW

IS C^3 . THAT'S TRUE BUT NONTRIVIAL.

SKETCH OF LOWER BOUND

$$u_\epsilon(x) \geq u(x) - C\epsilon$$

THIS TIME CONSIDER PAUL'S OPTIMAL v_k ,
AND A "TRIAL STRATEGY" FOR CAROL



$$x_{k+1} = x_k + \sqrt{2} \epsilon b_k v_k$$

UNTIL EXIT AT STEP K

$$u_\epsilon(x_k) = \max_{b \in \pm 1} u_\epsilon(x_k + \sqrt{2} \epsilon b_k v_k) + \epsilon^2$$

$$\geq u_\epsilon(x_k + \sqrt{2} \epsilon b_k v_k) + \epsilon^2$$

FOR ANY
STRATEGY
USED BY
CAROL

\Rightarrow $u_\epsilon(x_0) \geq K \epsilon^2$ WHERE EXIT OCCURS
AT STEP K

CLAIM: FOR GOOD CHOICE OF CAROL'S STRATEGY

$$u(x_{k+1}) \geq u(x_k) - \epsilon^2 + O(\epsilon^3)$$

IN FACT

$$\begin{aligned} \max_{b \in \pm 1} u(x_k + \sqrt{2} \epsilon b v) &= u(x_k) + \sqrt{2} \epsilon |v \cdot \nabla u| \\ &+ \epsilon^2 \langle D^2 u \cdot v, v \rangle + O(\epsilon^3) \end{aligned}$$

- IF $v \cdot \nabla u \neq 0$, THAT TERM DOMINATES
- IF $v \cdot \nabla u = 0$, THEN $v = \pm \nabla u^\perp / |\nabla u|$
- BOTH CASES GIVE DESIRED INEQ

So

$u(x_0) \leq K \epsilon^2 + O(\epsilon)$	WHERE EXIT OCCURS AT K
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THUS $u_\epsilon(x) \geq K \epsilon^2 \geq u(x_0) - C \epsilon$ QED

CONCLUSIONS

- NEW UNDERSTANDING (BUT NOT PHYSICAL?)
- SEMIDISCRETE NUMERICAL SCHEME (IS IT OF ANY PRACTICAL VALUE?)
- MANY OPEN QUESTIONS (WHICH 2nd ORDER EQUATIONS HAVE DETERMINISTIC CONTROL INTERPRETATIONS?)