

**Multiresolution Approximation of Curves with Normal
Polylines**

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Joint work with Ingrid Daubechies and Wim Sweldens.

Background

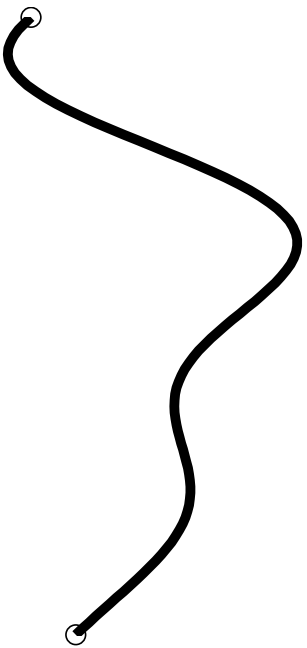
- Huge arbitrary topology meshes for real world 3D data.
- Goals:
 1. Compression,
 2. Progressive reconstruction,
 3. Signal processing, . . .
- How mesh is stored/represented affects compression rate.
Remeshing to better format improves rate.

Normal Meshes

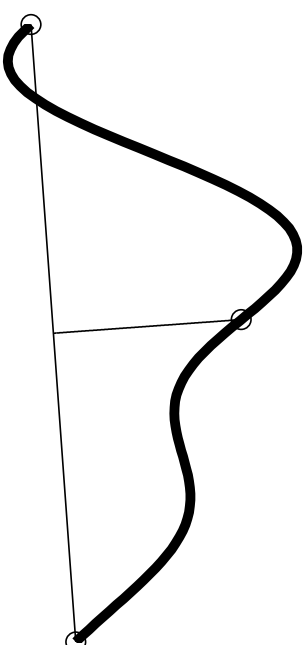
(Guskov, Khodakovsky, Schröder, Sweldens, . . .)

- Multiresolution meshes (predict + detail, at each level)
- Purely scalar representation (1 float/vertex),
- No “parameter” info
- Scalar wavelet compression codes can be used.

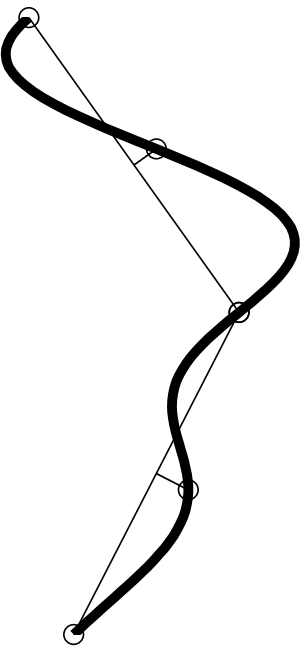
Construction of Normal Polyline



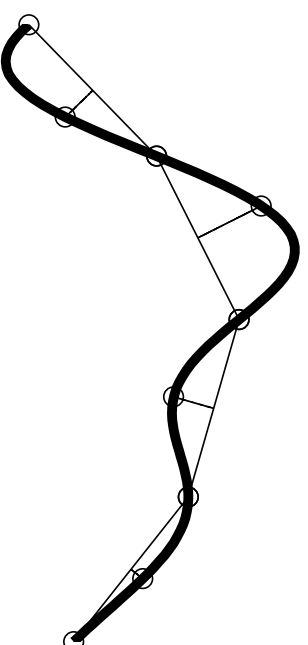
Starting curve $f = (x, y)^T$



Level 1



Level 2



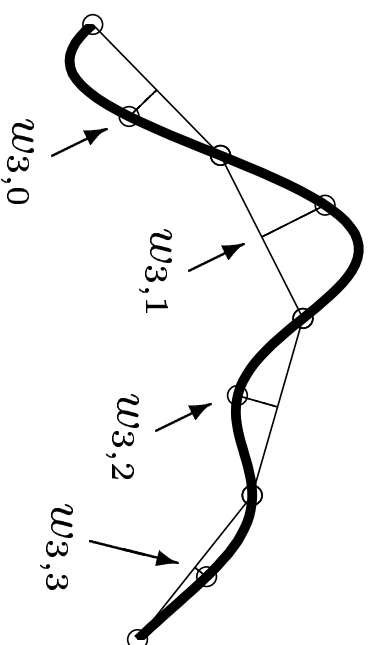
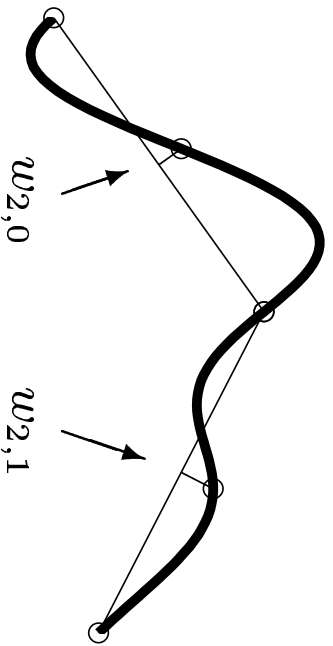
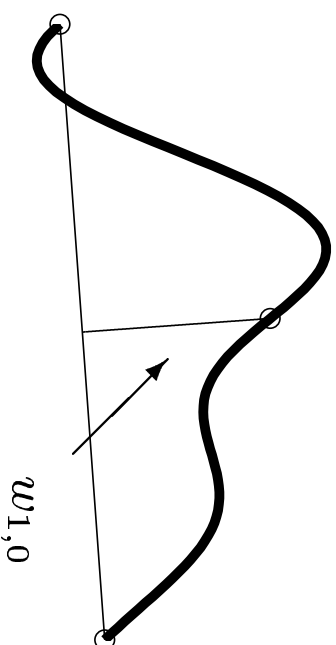
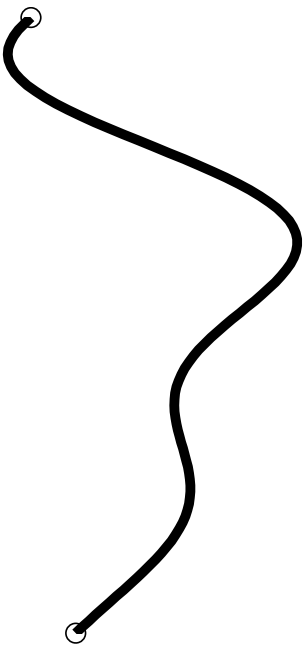
Level 3

One-dimensional analysis

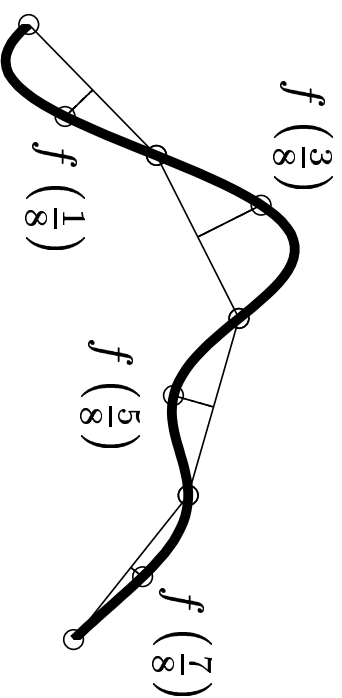
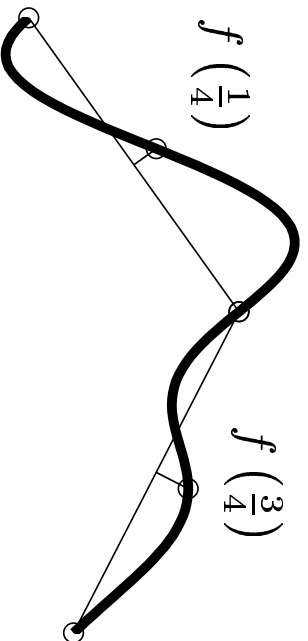
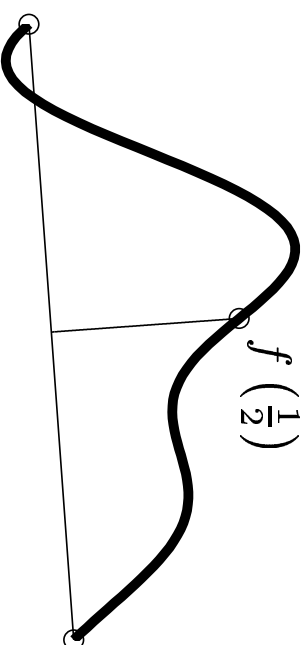
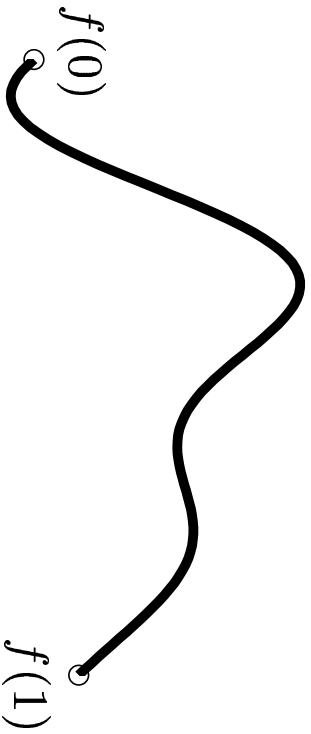
Issues:

- Decay of normal offsets (“wavelet coefficients”),
- Regularity of parameterization.

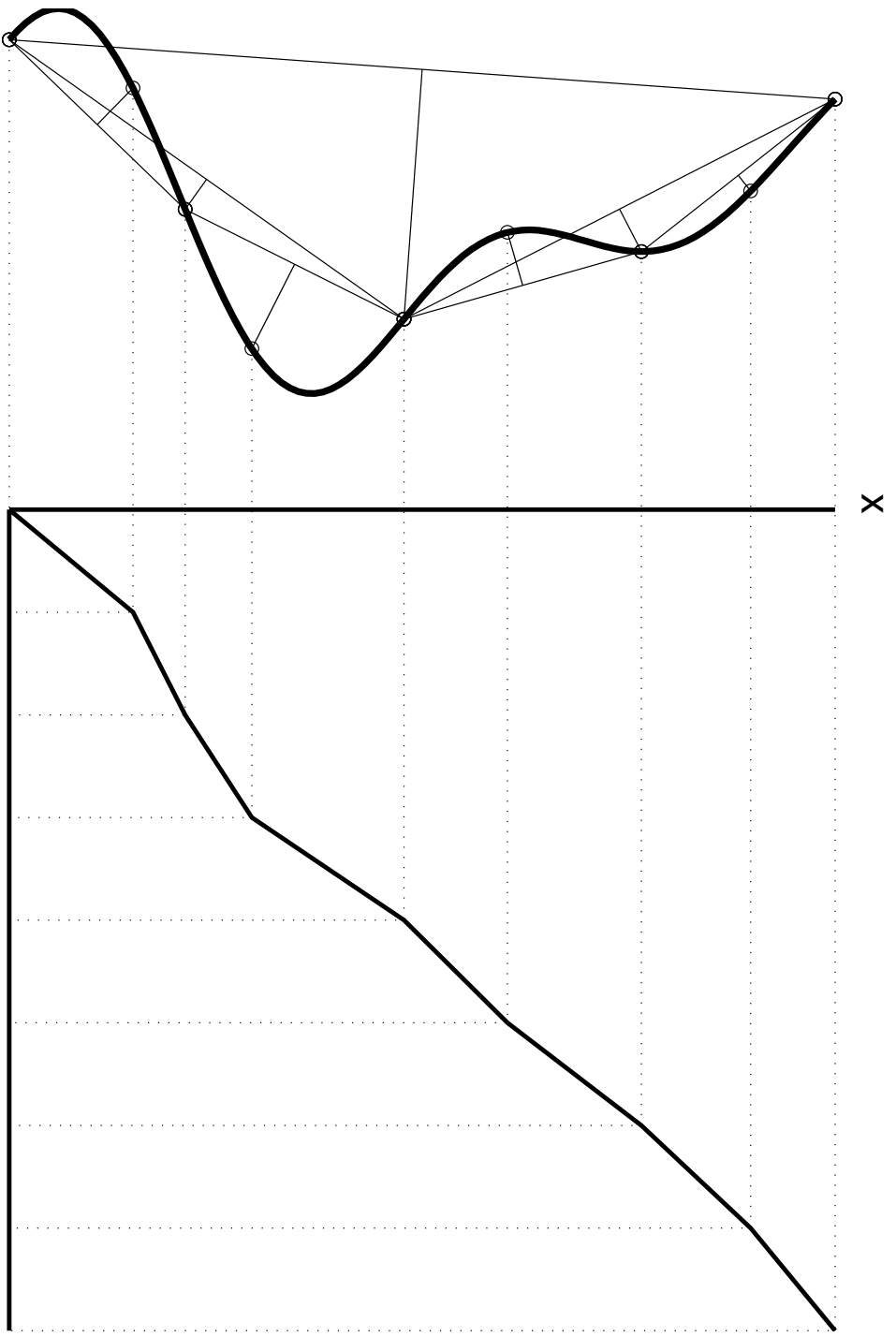
Normal Offsets $w_j = \{w_{j,k}\}$



Normal Parameterization $f(\alpha) = (x(\alpha), y(\alpha))^T$



Normal Parameterization, $x(\alpha)$



One-dimensional analysis, cont.

Parameterization typically *not* smooth even if curve is:

In linear case, let α be normal parameterization

$f(\alpha) = (x(\alpha), y(\alpha))^T$. Then

$$\left(f(\alpha + h) - f(\alpha - h) \right) \cdot \left(f(\alpha + h) - 2f(\alpha) + f(\alpha - h) \right) = 0,$$

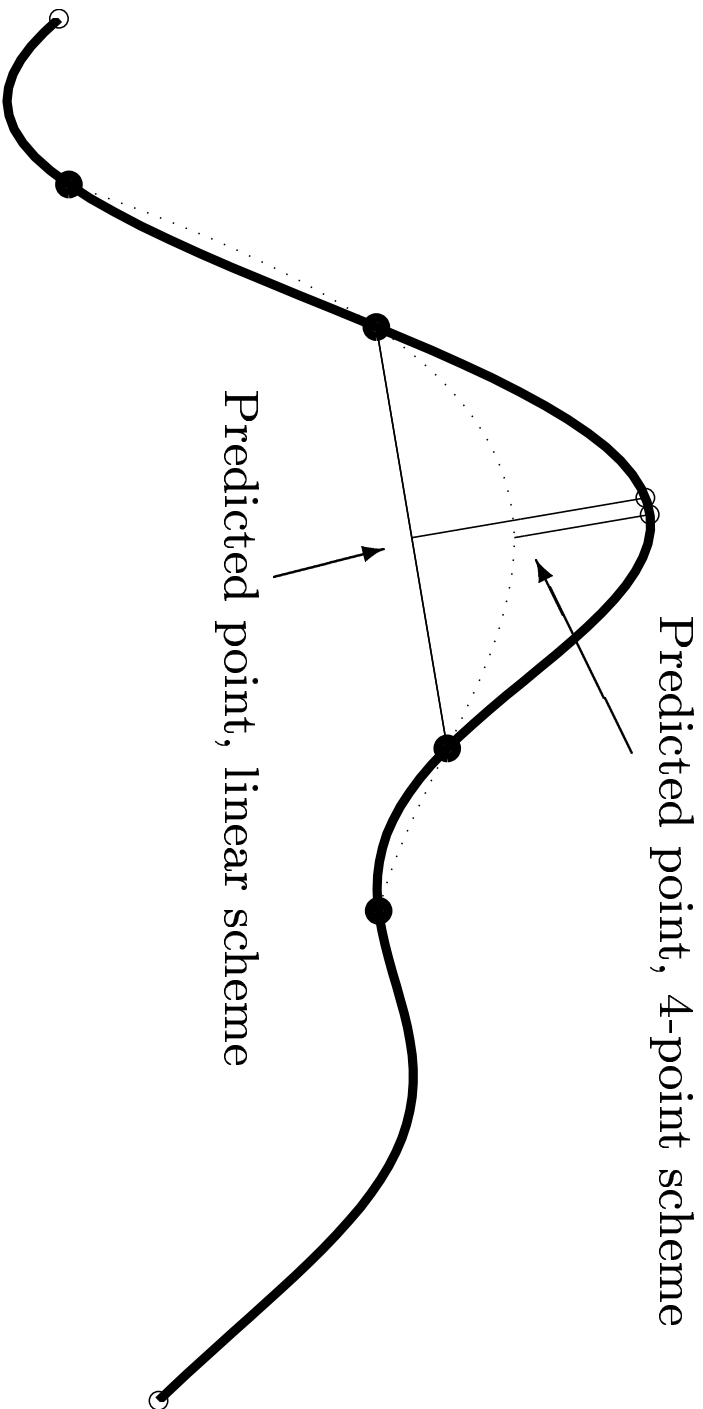
at odd dyadic points ($\alpha = (2k + 1)2^{-j}$, $h = 2^{-j}$).

Smooth parameterization \Rightarrow

$$\frac{d}{d\alpha} \left| \frac{df(\alpha)}{d\alpha} \right|^2 + \frac{h^2}{12} \left(\frac{d^3}{d\alpha^3} \left| \frac{df(\alpha)}{d\alpha} \right|^2 - \frac{d}{d\alpha} \left| \frac{d^2f(\alpha)}{d\alpha^2} \right|^2 \right) = \mathcal{O}(h^4).$$

$h \rightarrow 0 \Rightarrow$ parameterization \sim arclength and curve is line or circle.

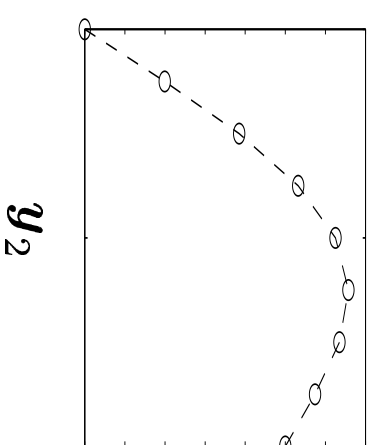
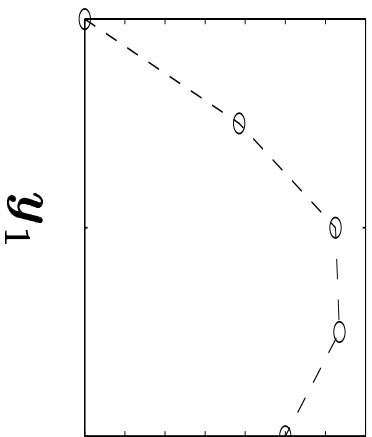
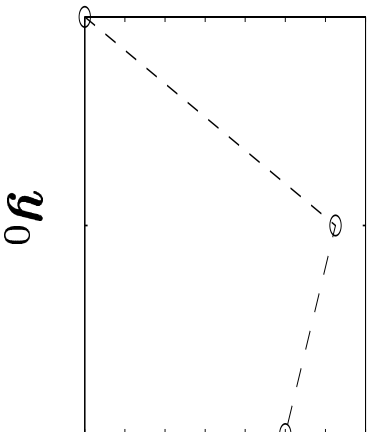
Higher Order Subdivision Scheme as Predictor



Subdivision

Procedure to generate successively finer mesh from a coarse mesh.

Ex. $(1D, \mathbf{y}_j = \{y_{j,k}\})$,



Let S be a *subdivision operator* such that $\mathbf{y}_{j+1} = S\mathbf{y}_j$ and

$$y_{j+1,k} = \sum_{\ell \in I_k} S_{k,\ell} y_{j,\ell}, \quad I_k = [k - B, k + B].$$

Here S is linear, stationary (indep. of j) and local ($B < \infty$).

Subdivision, cont.

Limit: $\lim_{j \rightarrow \infty} \tilde{\mathbf{y}}_j(t) \rightarrow \phi(t) \in C^r$

Order of S is p if S preserves q -degree polynomials for $0 \leq q < p$.

Derived subdivision schemes, $S^{[q]}$, act on divided differences of \mathbf{y}_j ,

$$S^{[q]} \mathbf{y}_j^{[q]} = \mathbf{y}_{j+1}^{[q]}, \quad \mathbf{y}_j^{[q]} = D_j^q \mathbf{y}_j.$$

(Also, $S^{[q]} = D_{j+1}^q S D_j^{-q}$.) $S^{[q]}$ well-defined for $q \leq$ order of S .

Examples, interpolating schemes ($y_{j+1,2k} = y_{j,k}$):

- “2-point” (linear), $S \sim [\frac{1}{2}, \frac{1}{2}]$, $\phi \in C^{1-\epsilon}$
- “4-point” (cubic), $S \sim \frac{1}{16}[-1, 9, 9, -1]$, $\phi \in C^{2-\epsilon}$.
- “6-point”, $S \sim \frac{1}{256}[3, -25, 150, 150, -25, 3]$, $\phi \in C^{2.83}$.

One-dimensional analysis, results

Theorem. Let S be predictor subdivision and $y(x)$ the curve. If

- S is of order $\geq p$ and, for some m , $\left|S^{[p]j}\right|_{\infty} \leq Cm^j$,
- $y \in C^{p+1}$,
- $|\Delta x_j|_{\infty} \leq C\delta^j$, for some $\delta < 1$.

Then

$$\tilde{x}_j(\alpha) \rightarrow x(\alpha) \in C^{p-\log_2 m-\varepsilon}.$$

and the normal offsets w_j satisfy the estimate

$$|w_j|_{\infty} \leq C2^{-j(p-\log_2 m+1-\varepsilon)}.$$

Examples:

- “2-point”, $S^{[1]} = I$, so $p = 1$, $m = 1$ and $x(\alpha) \in C^{1-\varepsilon}$,
- “4-point”, $\left|S^{[3]}\right|_{\infty} = 2$, so $p = 3$, $m = 2$ and $x(\alpha) \in C^{2-\varepsilon}$.

Sketch of proof

1. Consider nonlinear subdivision schemes T_j ,

$$\mathbf{x}_{j+1} = T_j \mathbf{x}_j, \quad |T_j \mathbf{x}_j - S \mathbf{x}_j|_\infty \leq 2^{-n_j}.$$

(a) Estimate $\mathbf{x}_j^{[q]}$, given $|S^{[p]^j}|_\infty \leq C m^j$,

$$\left| \mathbf{x}_j^{[q]} \right|_\infty \leq 1 + \begin{cases} (m 2^{q-p})^j, & n > p - \log_2 m, \\ (2^{q-n})^j, & n < p - \log_2 m. \end{cases}$$

(b) \Rightarrow regularity of limit,

$$\tilde{\mathbf{x}}(t) \rightarrow x(t) \in \begin{cases} C^{p - \log_2 m - \varepsilon}, & n > p - \log_2 m, \\ C^{n - \varepsilon}, & n < p - \log_2 m. \end{cases}$$

2. Show parameterization is given by some T_j with n large enough.

(a) Estimate (x -component of) offset,

$$|N_j \mathbf{x}_j - S \mathbf{x}_j|_\infty \leq |S y(\mathbf{x}_j) - y(S \mathbf{x}_j)|_\infty.$$

(b) Approximate commutation with nonlinear functions,

$$|\mathbf{x}^{[p]}|_\infty \leq \text{const} \Rightarrow |S y(\mathbf{x}) - y(S \mathbf{x})|_\infty \leq 2^{-j(p+1)}.$$

(c) Result follows by induction, starting from

$$|\mathbf{x}^{[1]}|_\infty \leq \text{const}$$

$$(2a, 2b) \Rightarrow |N_j \mathbf{x}_j - S \mathbf{x}_j|_\infty \leq |S y(\mathbf{x}_j) - y(S \mathbf{x}_j)|_\infty \leq 2^{-2j},$$

$$(1a) \Rightarrow |\mathbf{x}^{[2]}|_\infty \leq \text{const}$$

etc. till $|N_j \mathbf{x}_j - S \mathbf{x}_j|_\infty \leq 2^{-nj}$ for n large enough.