



One Size Fits All? : Computational Tradeoffs in Mixed Integer Programming Software

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A MIP Code is a Bag of Tricks

- ❑ **Presolve**
- ❑ **Cutting planes**
 - ❑ Gomory cuts
 - ❑ Knapsack cuts
 - ❑ Etc.
- ❑ **Node presolve**
- ❑ **Heuristics**
- ❑ **Node selection strategy**
- ❑ **Etc.**

Typical Path to Widespread Adoption of a New Technique

1. Improvement on a test set
2. Evaluation and tuning on a wider problem set
3. Implementation and deployment in a MIP code
4. Modelers (silently) benefit



A Few Examples

[Bixby, Fenelon, Gu, Rothberg, Wunderling, 2002]

<input type="checkbox"/> Cuts		53.7X
<input type="checkbox"/> Gomory	2.5X	
<input type="checkbox"/> MIR	1.8X	
<input type="checkbox"/> Knapsack	1.4X	
<input type="checkbox"/> Flow covers	1.2X	
<input type="checkbox"/> Implied bounds	1.2X	
<input type="checkbox"/> ...		
<input type="checkbox"/> Presolve		10.8X
<input type="checkbox"/> Heuristics		1.4X
<input type="checkbox"/> Node presolve		1.3X
<input type="checkbox"/> Probed dives		1.1X



Unlikely Path to Widespread Adoption of a New Technique

1. Improvement on a test set
2. Overall degradation on a wider problem set
3. Not implemented in a MIP Code
4. Modeler reads paper and implements technique



Another Unlikely Path to Widespread Adoption of a New Technique

1. **Improvement on a test set**
2. **Overall degradation on a wider problem set**
3. **Implemented in a MIP code anyway**
4. **Modeler:**
 1. Reads documentation or paper
 2. Recognizes that technique will be effective on his models
 3. Enables non-default option in MIP code



Example : Probing

[Brearley, Mitra, Williams, 1975]

- ❑ **Explore logical consequences of fixing binary variables to 0/1**
 - ❑ Variable fixing
 - ❑ Coefficient lifting
 - ❑ Implied bound cuts

- ❑ **Model mod011.mps**
 - ❑ Moderately difficult model from MIPLIB set



Model mod011.mps without probing

	Nodes		Objective	IInf	Best Integer	Cuts/ Best Node	ItCnt	Gap
Node	Left							
0	0		-6.2082e+007	16		-6.2082e+007	2254	
		...						
			-5.8833e+007	48		Flowcuts: 5	6840	
*	52	52		0	-4.1232e+007	-5.8811e+007	12548	42.64%
*	53	50		0	-5.1414e+007	-5.8811e+007	12565	14.39%
*	156	114		0	-5.3439e+007	-5.8563e+007	27821	9.59%
*	319	242		0	-5.3676e+007	-5.8305e+007	52401	8.62%
*	955	716		0	-5.3779e+007	-5.7598e+007	180050	7.10%
*	976	721		0	-5.3833e+007	-5.7545e+007	183880	6.89%
*	1168	858		0	-5.3863e+007	-5.7390e+007	217373	6.55%
*	1187	868		0	-5.3898e+007	-5.7388e+007	220889	6.47%
*	1748	1205		0	-5.4059e+007	-5.7096e+007	331173	5.62%
*	2393	1601		0	-5.4088e+007	-5.6805e+007	475812	5.02%
*	2569	1395		0	-5.4422e+007	-5.6752e+007	511588	4.28%
*	12152	3545		0	-5.4559e+007	-5.5205e+007	2175100	1.18%

Implied bound cuts applied: 410

Flow cuts applied: 695

Flow path cuts applied: 117

Gomory fractional cuts applied: 14

Integer optimal solution: Objective = -5.4558535014e+007

Solution time = 2572.66 sec. Iterations = 2444566 Nodes = 16437



Model mod011.mps with probing

Reduced MIP has 1558 rows, 6895 columns, and 14668 nonzeros.

...

Probing added 286 nonzeros
Probing time = 0.31 sec.

	Nodes				Cuts/			
	Node	Left	Objective	IInf	Best Integer	Best Node	ItCnt	Gap
	0	0	-5.9923e+007	16		-5.9923e+007	2165	
			...					
			-5.4636e+007	33		Cuts: 10	25098	
*	33	33		0	-5.4195e+007	-5.4636e+007	34600	0.81%
*	51	38		0	-5.4219e+007	-5.4634e+007	43147	0.77%
*	78	33		0	-5.4559e+007	-5.4625e+007	57046	0.12%

Implied bound cuts applied: 2398
Flow cuts applied: 496
Flow path cuts applied: 207
Gomory fractional cuts applied: 3

Integer optimal solution: Objective = -5.4558535014e+007
Solution time = 206.73 sec. Iterations = 59222 Nodes = 110



Probing on a Wider Set of Models

□ Mean performance ratio:

□ 1s 0.68

□ 10s 0.59

□ 100s 0.34

□ 500s 0.31

□ 1000s 0.25



Another Example : Strong Branching

[Applegate, Bixby, Chvatal, and Cook, 1995]

- **Use dual simplex to choose branching variable**
 - Estimate objective degradation by performing a limited number of simplex iterations
 - Maximize the minimum degradation
- **Finding optimal solutions for large Traveling Salesman Problems**
 - Crucial for improving objective lower bound



Strong Branching on a Wider Set of Models

□ Mean performance ratio:

□ 1s	0.96
□ 10s	0.76
□ 100s	0.72
□ 500s	0.66
□ 1000s	0.69



Consistency Between User Goals and Code Goals



A MIP Code Has An (Implicit) Emphasis

- ❑ **Emphasis before CPLEX 7.0:**
 - ❑ Minimize time to proven optimality
- ❑ **Important components of approach:**
 - ❑ Aggressive cut generation
 - ❑ Search strategy that attempts to avoid unnecessary work
 - Depth First Search until first feasible found
 - Best Bound Search until termination

Potential Mismatch Between Goals

❑ User reaction

- ❑ Depends heavily on user metric
- ❑ Common metric:
 - Time to first “good” feasible solution

❑ Reevaluate the bag of tricks

- ❑ Time to first feasible?
- ❑ Time to 10% (20%?) gap?



Performance for Feasibility Emphasis

□ Mean performance improvement (CPLEX 7.5):

	Desired optimality gap			
	10%	20%	30%	finite
□ 1s	1.00	0.99	1.00	1.01
□ 10s	1.08	1.10	1.12	1.25
□ 100s	1.17	1.26	1.31	1.51
□ 500s	1.08	1.05	1.24	1.50
□ 1000s	1.27	1.40	1.47	1.46

User Emphasis Setting: Feasibility Instead of Optimality

- ❑ **Simple underlying algorithmic changes**
 - ❑ Less aggressive application of cuts
 - ❑ More time spent near leaves of search tree
- ❑ **Could be achieved with parameter changes**
- ❑ **Users pleased nonetheless**
 - ❑ Describe goals rather than understanding and choosing techniques

User Emphasis Setting in 8.0

- ❑ **Improvements reduce the importance of emphasis:**
 - ❑ More heuristics produce feasible solutions faster and more consistently
 - ❑ Probed dives make dives more likely to lead to good feasible solutions

Performance for Feasibility Emphasis

□ Mean performance improvement (CPLEX 8.0):

	Desired optimality gap			
	10%	20%	30%	Finite
□ 1s	1.03	1.02	1.02	1.01
□ 10s	1.02	1.01	1.02	0.99
□ 100s	1.16	1.14	1.13	1.04
□ 500s	1.08	1.09	1.03	1.08
□ 1000s	0.72	0.79	0.77	0.97

MIP Results of CPLEX8.0

[Bixby, Fenelon, Gu, Rothberg, Wunderling, 2002]

- ❑ **Test set: 978 models**
 - Selected from our library with over 1500 models**
- ❑ **100,000 seconds limit on ES40 Compaq Alpha**
- ❑ **Solved to optimality**
775 (77%)
- ❑ **Among those not solved to optimality**
 - 116 had gap less than 10% (11.9%)**
 - 32 had no integral solution (3.2%)**
- ❑ **MIP emphasis feasibility on the 32 models**
 - 25 found no feasible solution (2.6%)**



Hard Problems

- **Natural progression:**
 1. Try default settings
 2. Specify an emphasis
 3. Change parameter settings
 4. Use priority order
 5. Reformulate model
- **Some models still unsolved after all these steps**





Exploiting User Knowledge



User Knowledge

- ❑ **Users sometimes have domain knowledge that can help solution**
 - ❑ Crucial variables
 - ❑ Heuristics for finding good feasible solutions
 - ❑ Strategies to decompose the problem through branching
 - ❑ Special cutting planes
 - ❑ Etc.
- ❑ **User knowledge on the original model**
 - ❑ MIP code solves the *presolved* model



Advanced Features

❑ MIP callbacks

- ❑ Cut callback
- ❑ Heuristic callback
- ❑ Branch callback
- ❑ Incumbent callback

❑ Advanced presolve

- ❑ Allows user to express domain knowledge in terms of the original model

Adding User Cuts or Lazy Constraints

❑ Before CPLEX 7.0

- ❑ Usually need to turn off presolve (lose 10.8X)
- ❑ All constraints must be explicit

❑ Since CPLEX 7.0

- ❑ Can choose whether callbacks will work with original or presolved model
- ❑ Can obtain mappings for variables and constraints in the original model
- ❑ No need to specify all constraints up front
 - *Lazy constraints*



Caution: User Cuts

- ❑ **Original model**

$$\max \{x_1 + x_2 + 2x_3 : 3x_1 + 3x_2 + 4x_3 \leq 6, x_1, x_2, x_3 \in \mathbb{B}\}$$

- ❑ **Presolved model**

$$\max \{y + 2x_3 : 3y + 4x_3 \leq 6, 0 \leq y \leq 2, 0 \leq x_3 \leq 1, y, x_3 \in \mathbb{Z}\}$$

- ❑ **Cut for the original model**

$$x_1 + x_3 \leq 1$$

It cannot be transformed and added to the presolved model

- ❑ **Need to turn off non-linear reductions, such as parallel column reduction**



Caution: Lazy Constraints

□ Model

$$\max \{x: 5x + 3y \leq 10, x - y \leq 0, x \geq 0, y \geq 0, x \in \mathbb{Z}\}$$

□ $x - y \leq 0$ treated as a lazy constraint

Presolve will fix y to 0 and x to 2

$x - y \leq 0$ becomes $2 - 0 \leq 0$

Augmented model is infeasible?

□ Need to turn off *dual reductions*

□ Reductions that depend on the objective function





Side Constraints



Side Constraints

- ❑ **Many types of “side constraints”**
 - ❑ SOS constraints
 - ❑ Semi-continuous variables
 - ❑ Cardinality constraints
 - ❑ Min, Max and Abs functions
 - ❑ Logical expressions
 - e.g., $x = 1$ implies $y+z \leq 3$
 - ❑ Tour requirements (TSP)
 - ❑ Etc.



Handling Side Constraints - Linearize

□ Example: SOS1(z_1, z_2)

- Introduce auxiliary binary variables b_1, b_2

- $z_1 \leq u_1 b_1; \quad z_2 \leq u_2 b_2; \quad b_1 + b_2 \leq 1$

□ Pros:

- All MIP tricks apply (cuts, presolve, heuristics, etc.)

- No need to handle special cases in MIP code

□ Cons:

- Model size increases

- Often leads to large big-M coefficients



Handling Side Constraints - Branching

□ Example: SOS1(z_1, z_2)

- When both $z_1 > 0$ and $z_2 > 0$ at a node...
- Branch on SOS1:
 - Left child: $z_1 = 0$
 - Right child: $z_2 = 0$

□ Cons:

- Special case for each construct
- No presolve, cuts, heuristics, ...
- Looser relaxation



Tighter Implicit Formulation?

- ❑ Is it possible to tighten relaxation without an explicit linearization?
- ❑ **Specialized cuts or lazy constraints**
 - ❑ E.g., cardinality constraints [de Farias and Nemhauser], TSP [Applegate, Bixby, Chvátal, and Cook]
 - ❑ Need to derive for each type of non-linear constraint
- ❑ **Alternative: extension to Gomory cuts**

Gomory Cut Review

□ Given $y, x_j \in Z_+$, and

$$y + \sum a_{ij}x_j = d = \lfloor d \rfloor + f, \quad f > 0$$

□ **Rounding:** Where $a_{ij} = \lfloor a_{ij} \rfloor + f_j$, define

$$t = y + \sum (\lfloor a_{ij} \rfloor x_j : f_j \leq f) + \sum (\lceil a_{ij} \rceil x_j : f_j > f) \in Z$$

□ **Then**

$$\sum (f_j x_j : f_j \leq f) + \sum (f_j - 1) x_j : f_j > f = d - t$$

□ **Disjunction:**

$$t \leq \lfloor d \rfloor \Rightarrow \sum (f_j x_j : f_j \leq f) \geq f$$

$$t \geq \lceil d \rceil \Rightarrow \sum ((1 - f_j) x_j : f_j > f) \geq 1 - f$$

□ **Combining:**

$$\sum ((f_j / f) x_j : f_j \leq f) + \sum (((1 - f_j) / (1 - f)) x_j : f_j > f) \geq 1$$



An Important Class: Disjunctive Constraints

□ Typical disjunctive set of constraints

x must satisfy at least k of n sets of linear constraints, $S_i = \{x: A_i x \geq b_i\}$ for $i = 1, \dots, n$

□ Modeling with binary variables

Dantzig (1957) , Nemhauser and Wolsey (1988)

□ Side constraints in the class

- SOS constraints
- Semi-continuous variables
- Cardinality constraints
- Min, Max and Abs functions
- Logical linear expressions



Cardinality Constraint

□ Definition

At most m variables of x_1, \dots, x_n can be positive

□ Use typical disjunctive set to express

$$S_i = \{x: -x_i \geq 0\} \text{ for } i = 1, \dots, n$$

$$k = n - m$$



Gomory Cut Extension (with Puget)

- **Given** $x_j \in R_+$, and

x is not in S_1, S_2, \dots, S_m , with $m > n - k$

Note x should be in at least $k - (n - m)$ of the above sets

- **Pick a violated constraints from each set**

$$\sum a_{ij} x_j \geq d_i, i = 1, \dots, m$$

- **Substitute basic variables with nonbasic ones**

$$\sum f_{ij} x_j \geq g_i, i = 1, \dots, m$$

Note $g_i > 0$. Let $h_{ij} = \max(0, f_{ij} / g_i)$, then

$$\sum h_{ij} x_j \geq 1, i = 1, \dots, m$$

- **Combine**

$$\sum \sum h_{ij} x_j \geq m + k - n$$



One Size Fits All?

- ❑ **Default works well to prove optimality or to find good feasible solutions for most models**
 - ❑ Try it first
- ❑ **CPLEX has an emphasis setting.**
 - ❑ Using it to specify a goal may help for some models
- ❑ **Several features are off by default**
 - ❑ Turning them on or changing parameter settings can help solving hard models
- ❑ **CPLEX provides advanced routines for exploiting user knowledge**
 - ❑ They can be helpful for hard models, e.g. their use for extending Gomory cuts for handling side constraints